

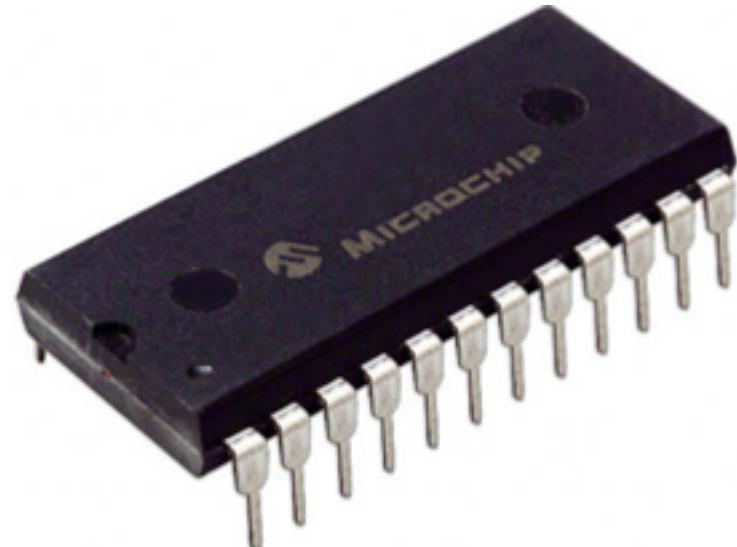
Robustly Reusable Fuzzy Extractor from Standard Assumptions



Yunhua Wen and Shengli Liu
Shanghai Jiao Tong University

Problem

- **Randomness** is crucial in cryptography (e.g. sk).
- However, **uniformly distributed and accurately reproducible** string is rare in practice.
- There are many imperfect random sources, e.g.



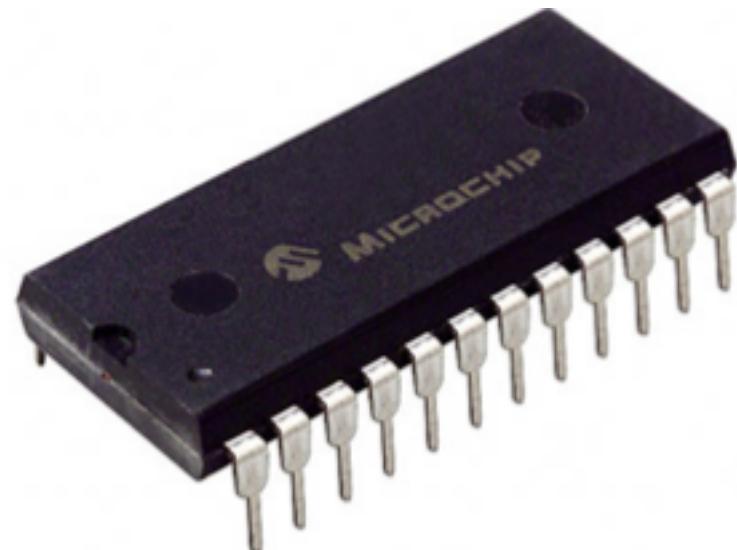
Physically Unclonable Functions (PUFs)



Biometric Information

Problem

- **Randomness** is crucial in cryptography (e.g. sk).
- However, **uniformly distributed and accurately reproducible** string is rare in practice.
- There are many imperfect random sources, e.g.



Physically Unclonable Functions (PUFs)



Biometric Information

Problem: How to use such imperfect random sources in cryptography?

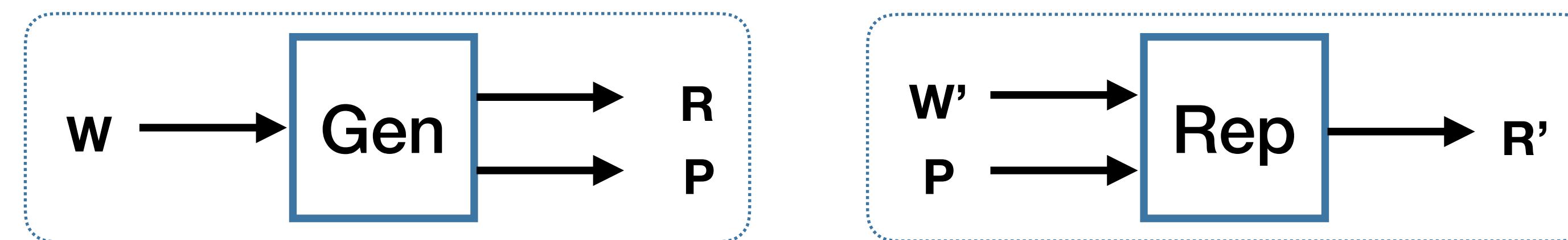
Fuzzy Extractor

- **Gen(w)**

- **Input:** a weak random secret w .
- **Output:** the extracted key R and a public helper string P .

- **Rep(w' , P)**

- **Input:** a noisy version w' and the public helper string P .
- **Output:** the extracted key R' .



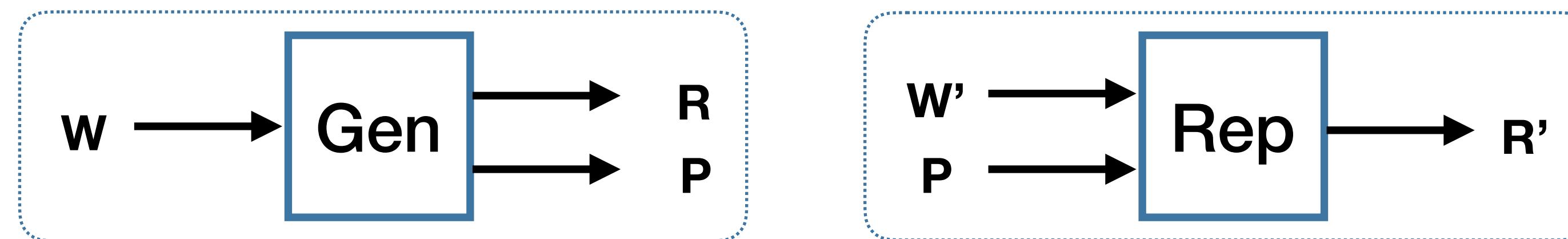
Fuzzy Extractor

- **Gen(w)**

- **Input:** a weak random secret w .
- **Output:** the extracted key R and a public helper string P .

- **Rep(w' , P)**

- **Input:** a noisy version w' and the public helper string P .
- **Output:** the extracted key R' .

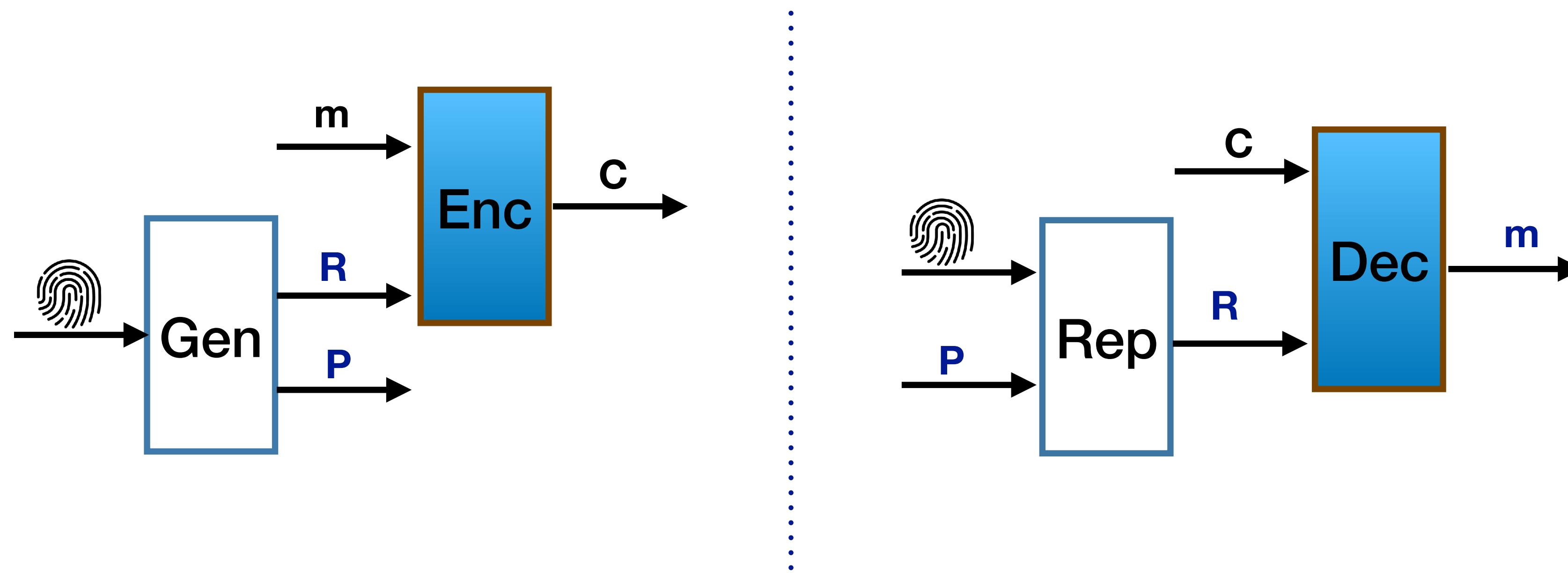


Correctness: If w' is close enough to w , $R' = R$.

Security: R is pseudorandom given P .

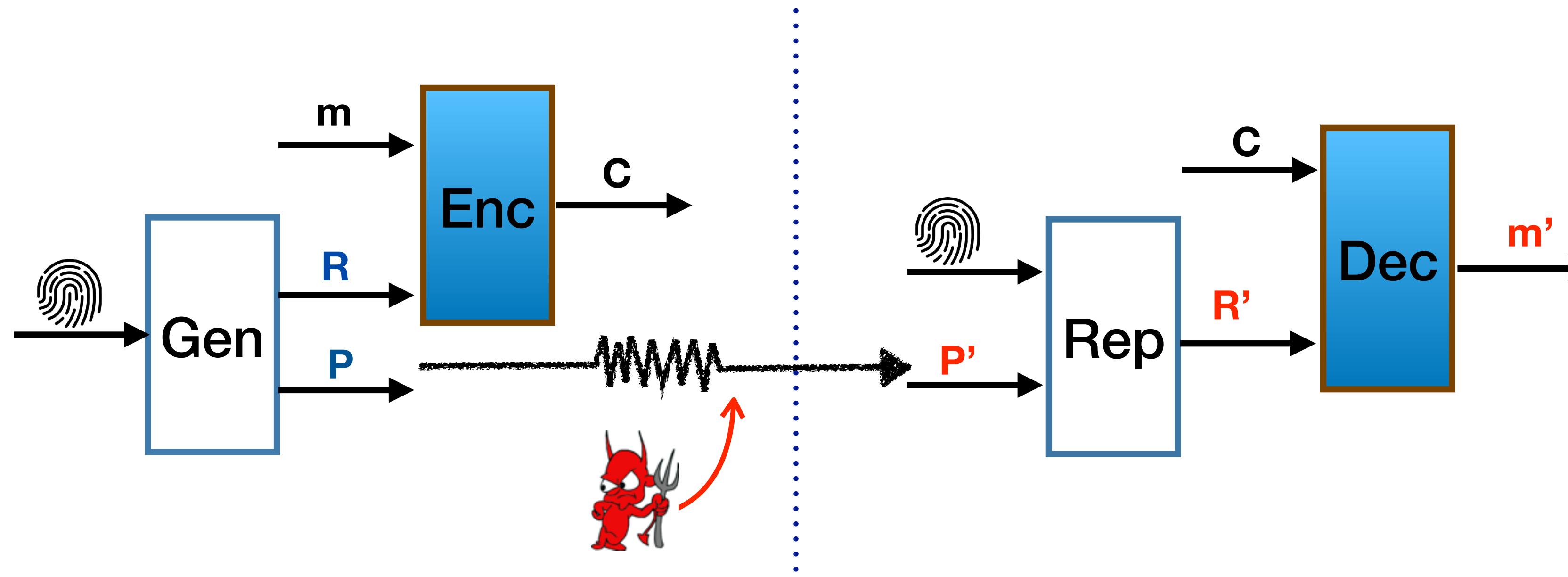
Applications

Application in **Encryption and Decryption**:



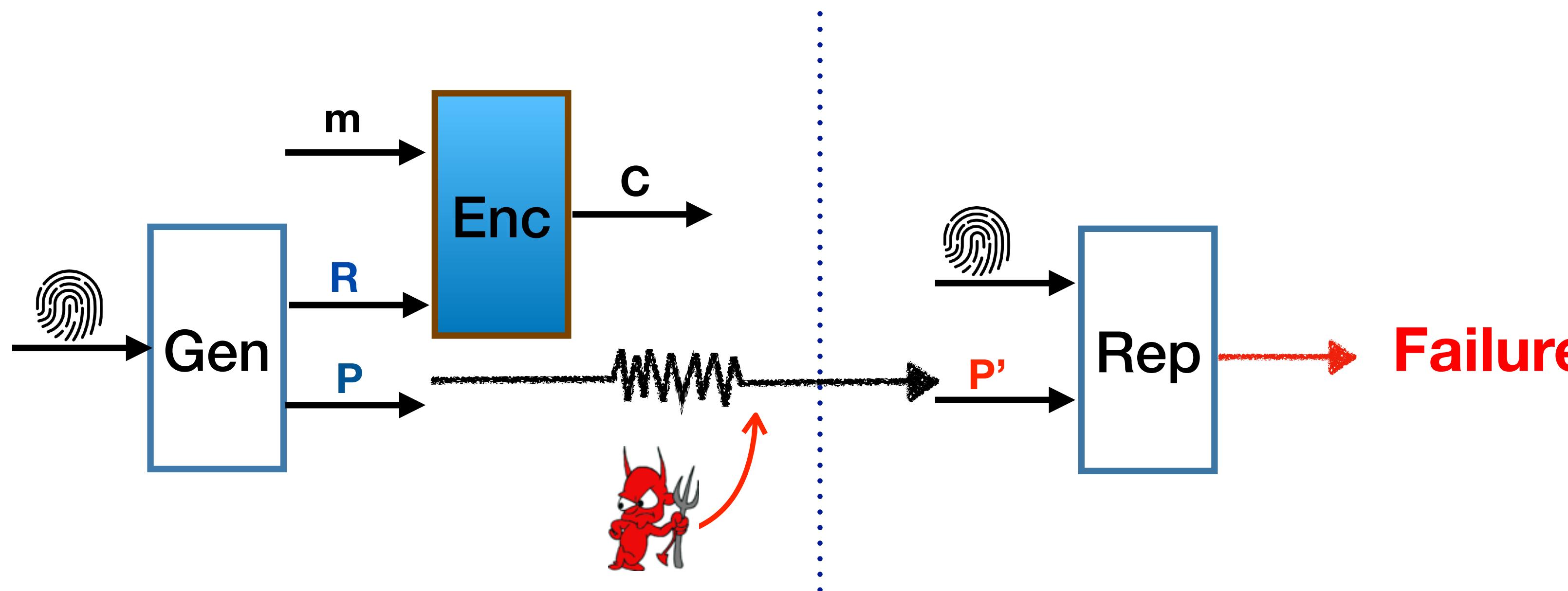
Users do not need to store the secret key **R**.

Robust Fuzzy Extractor



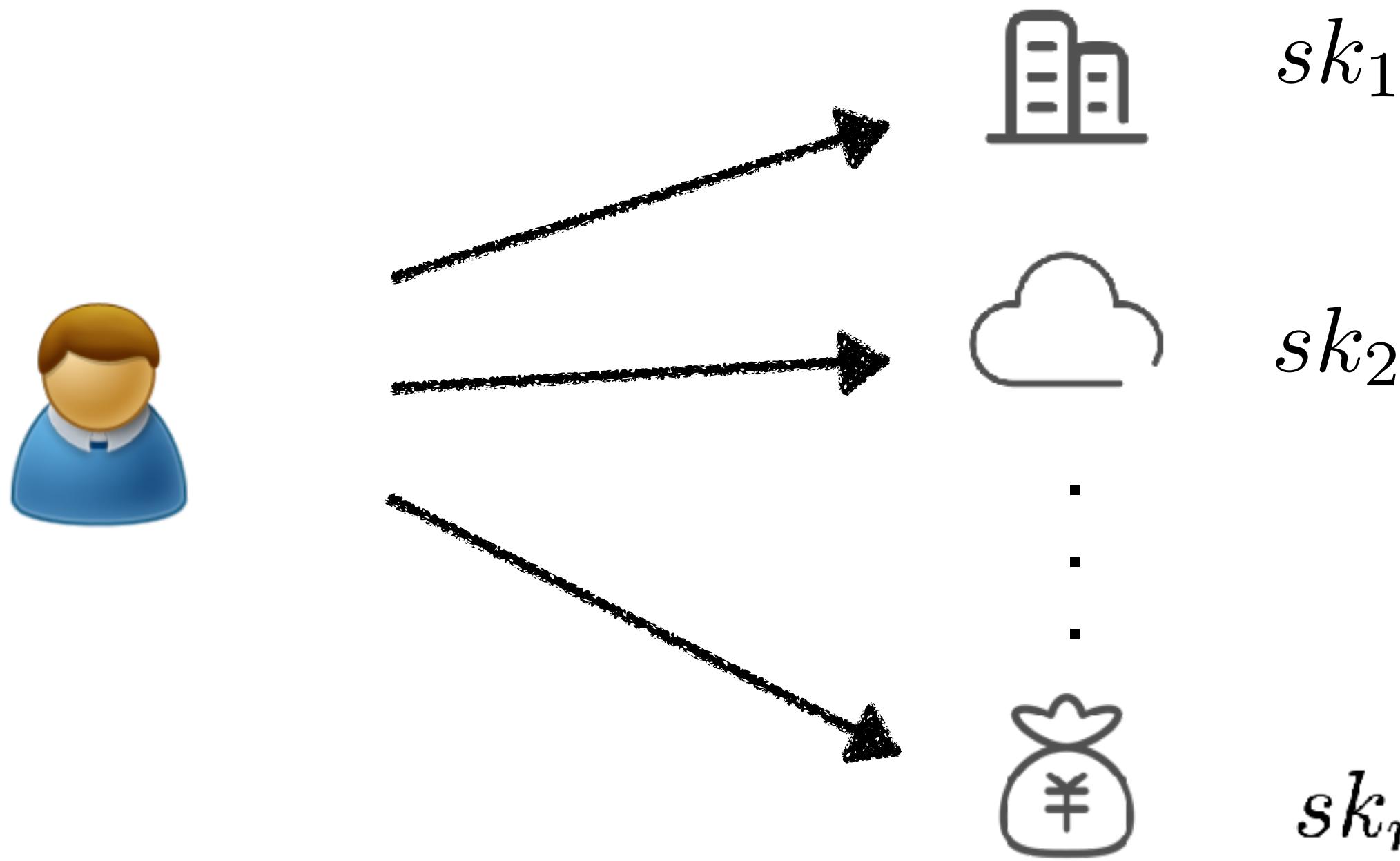
The user may get a **wrong** key **R'** without notifications.

Robust Fuzzy Extractor



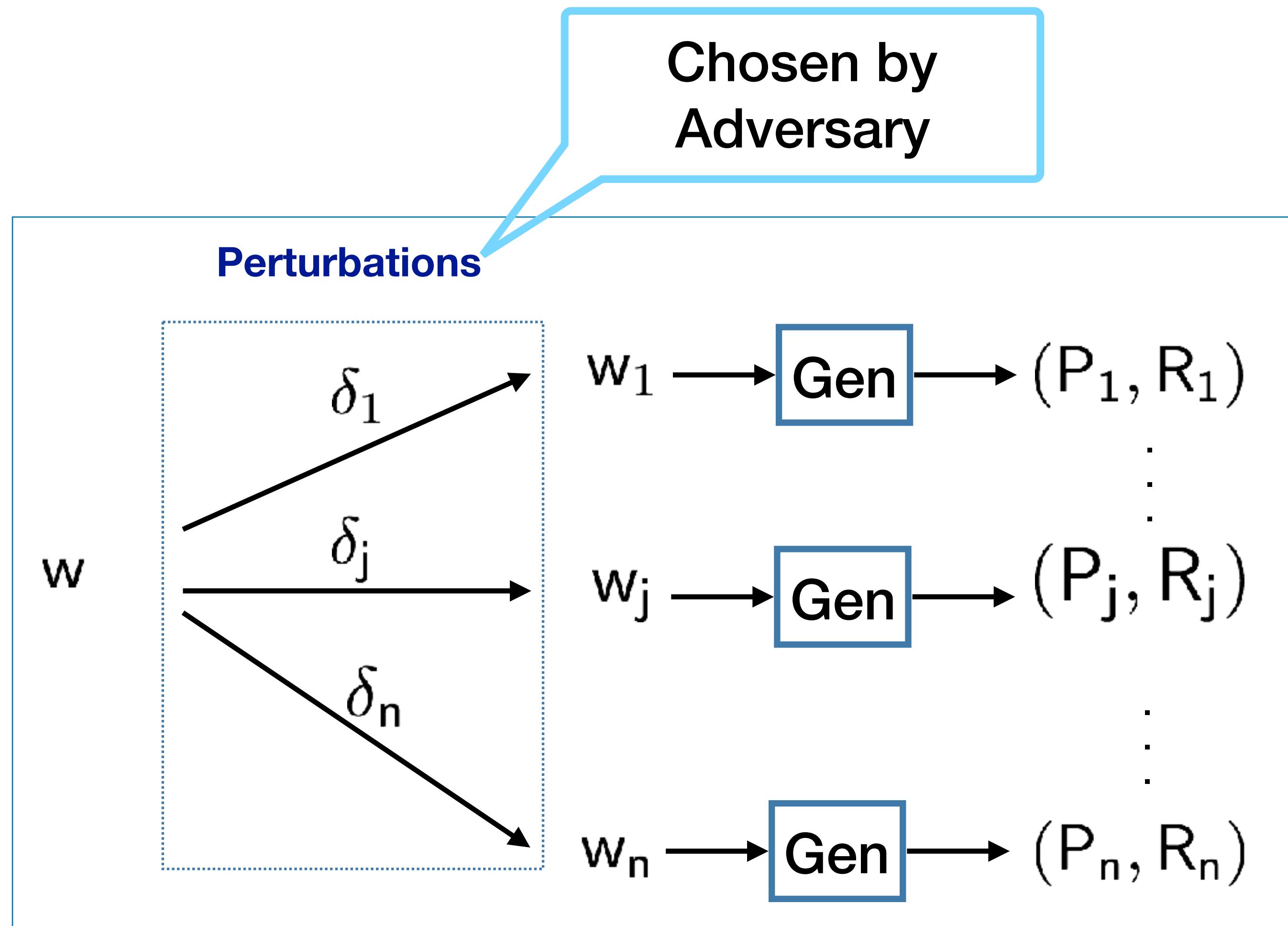
Security: If P is modified, then Rep will output \perp .

Reusable Fuzzy Extractor

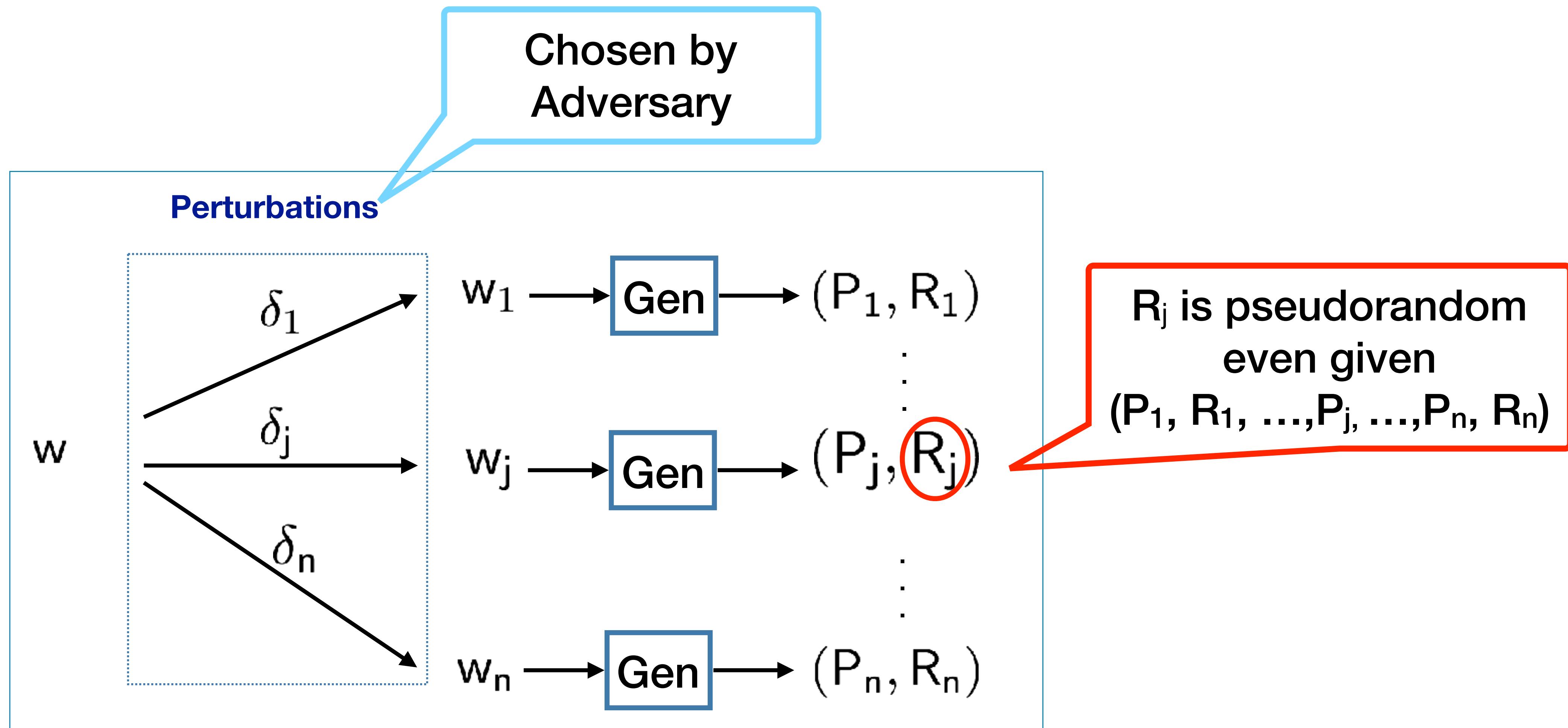


- Biometric is unique and cannot be **changed** or **created**.
- The security of **multi-extraction** from the **same noisy source** is not guaranteed by fuzzy extractor.

Reusable Fuzzy Extractor



Reusable Fuzzy Extractor



Related Works

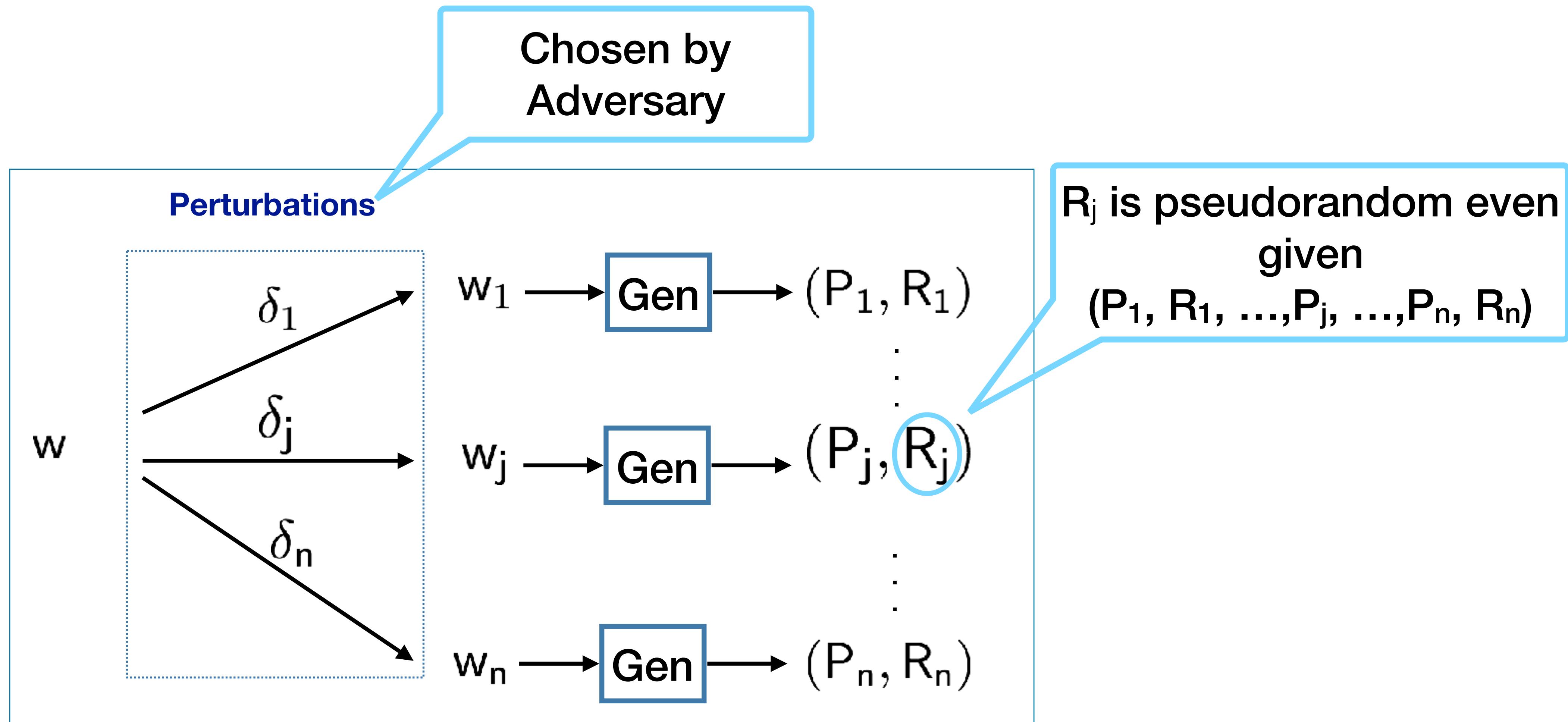
FE schemes	Robustness?	Reusability?
[DRS04], [FMR13]	✗	✗
[Boyen04], [ABCG16], [CFPRS16], [ACEK17], [WL18], [WLH18]	✗	✓
[BDKOS05], [DKRS06], [KR08], [CDFPW08]	✓	✗

No fuzzy extractor considers robustness and reusability simultaneously.

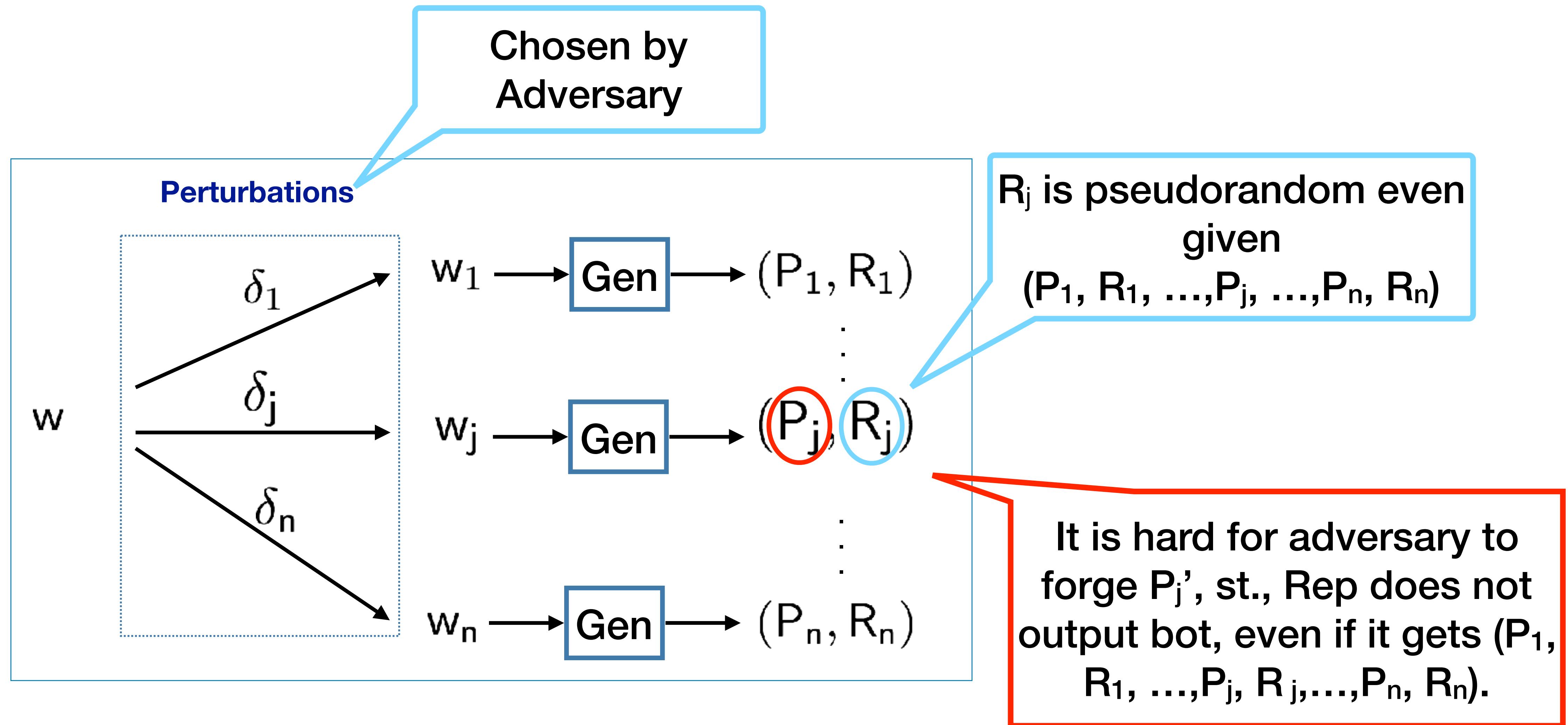
Our Contribution

- We formally defined robustly reusable fuzzy extractor(rrFE).
- We constructed the first rrFE based on standard assumptions.

Robustly Reusable Fuzzy Extractor



Robustly Reusable Fuzzy Extractor

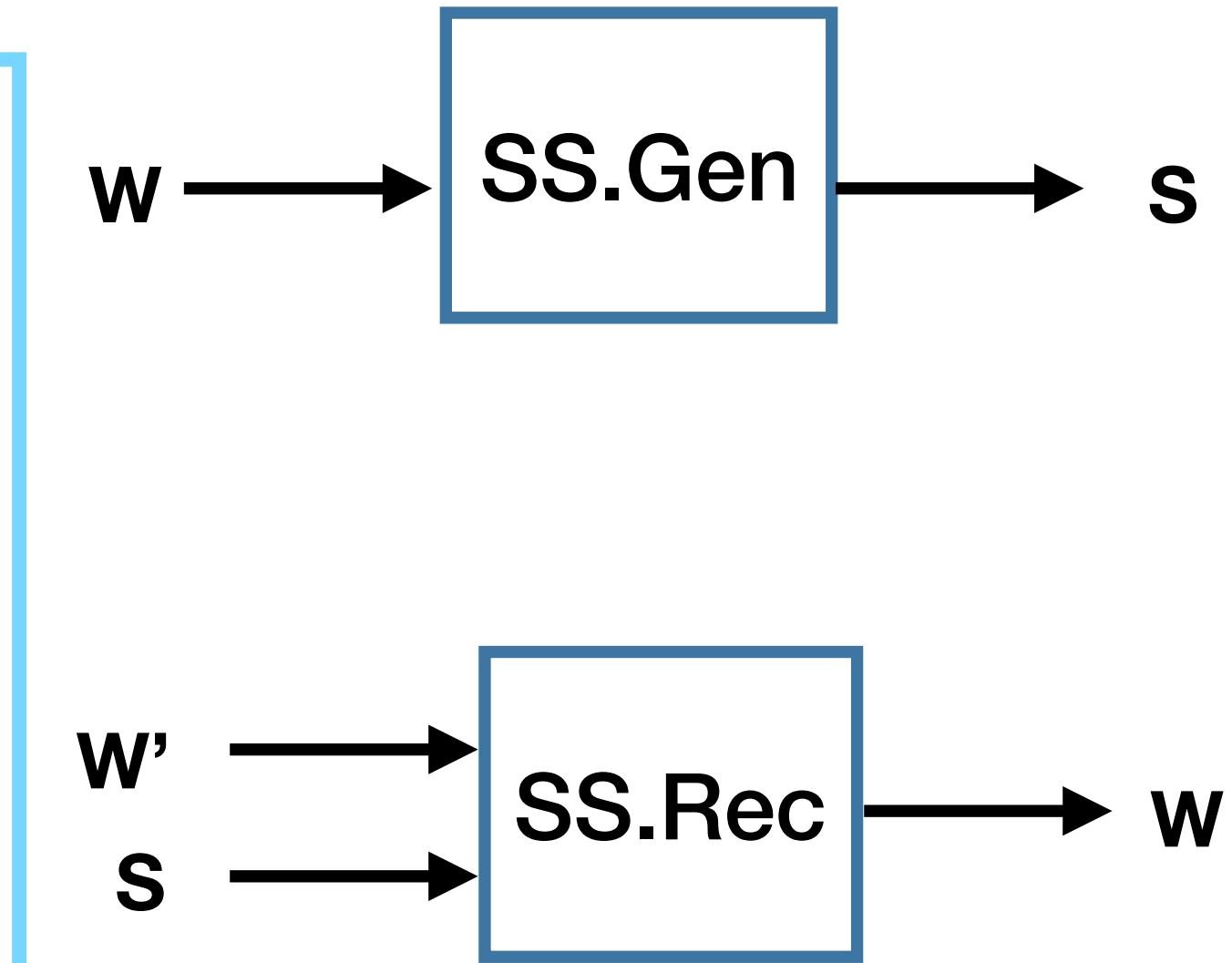


Building Blocks

- Homomorphic Secure Sketch (SS)
- Homomorphic Extractor (Ext)
- Symmetric Key Encapsulation Mechanism (SKEM)
- Homomorphic Lossy Algebraic Filter (LAF)

Building Block-Secure Sketch

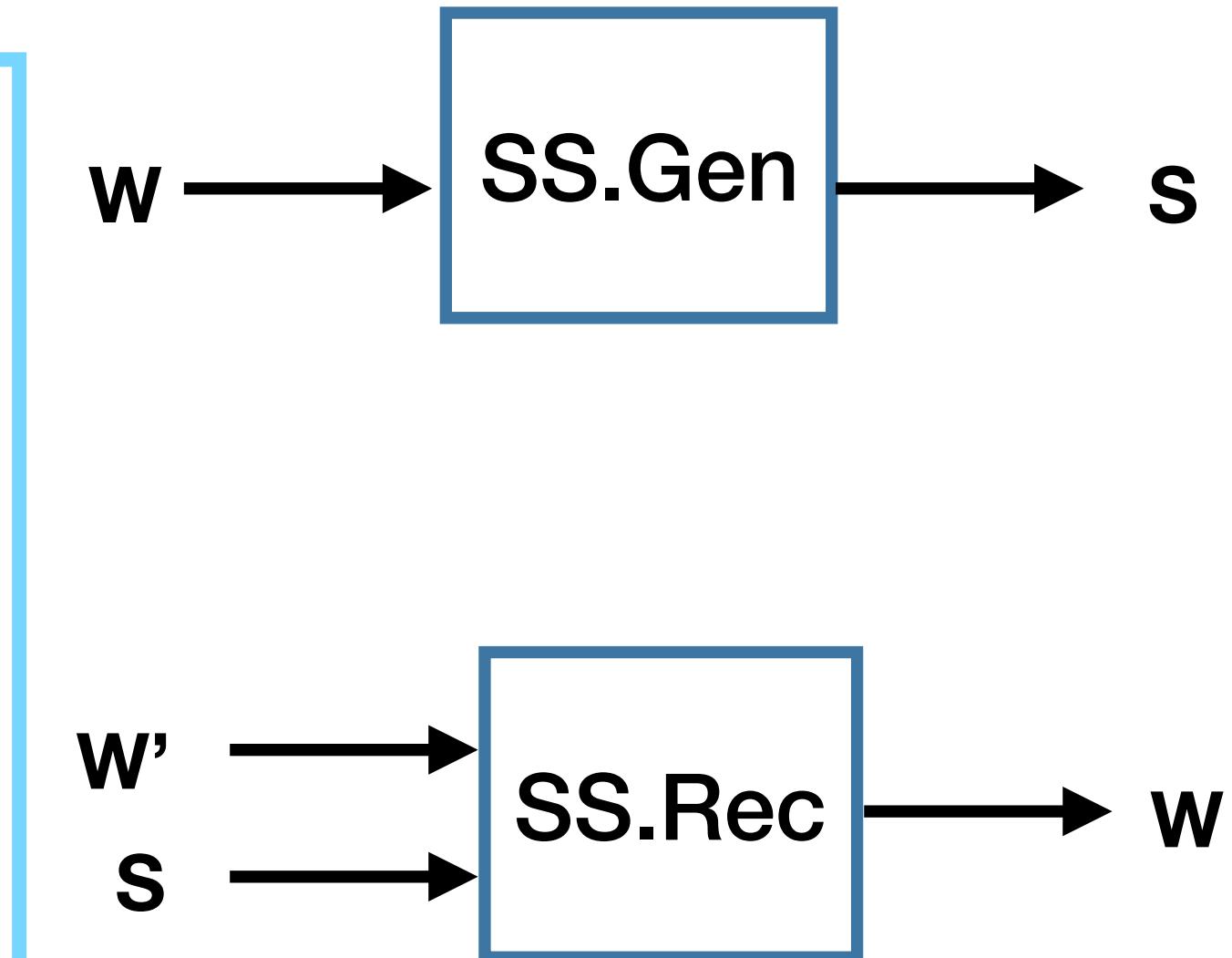
- **SS.Gen(w)**
 - Input: a weak random secret w .
 - Output: a sketch s .
- **SS.Rec(w' , s)**
 - Input: a noisy version w' and the sketch s .
 - Output: w .



- **Correctness:** For w' close to w , w can be recovered from s .
- **Privacy:** s does not leak too much information of w .

Building Block-Secure Sketch

- **SS.Gen(w)**
 - Input: a weak random secret w .
 - Output: a sketch s .
- **SS.Rec(w' , s)**
 - Input: a noisy version w' and the sketch s .
 - Output: w .



- **Correctness:** For w' close to w , w can be recovered from s .
- **Privacy:** s does not leak too much information of w .

Homomorphic secure sketch:

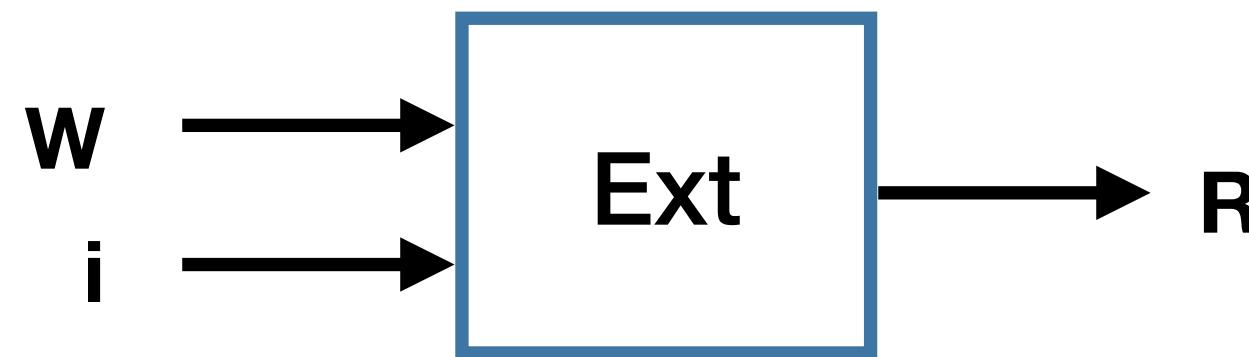
$$\text{SS.Gen}(w+w') = \text{SS.Gen}(w) + \text{SS.Gen}(w').$$

Building Block – Extractor

Extractor

Input: a weak secret w and a uniformly random seed i .

Output: extracted key $R = \text{Ext}(w; i)$.



Security: R is uniformly random, even conditioned on the seed i .

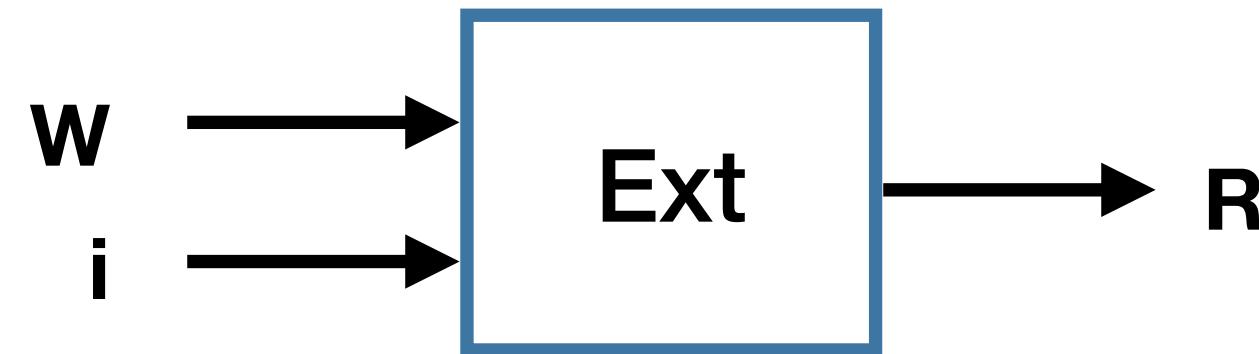
$$(\text{Ext}(W; i), i) \approx (\text{Uniform}, i).$$

Building Block – Extractor

Extractor

Input: a weak secret w and a uniformly random seed i .

Output: extracted key $R = \text{Ext}(w; i)$.



Security: R is uniformly random, even conditioned on the seed i .

$$(\text{Ext}(W; i), i) \approx (\text{Uniform}, i).$$

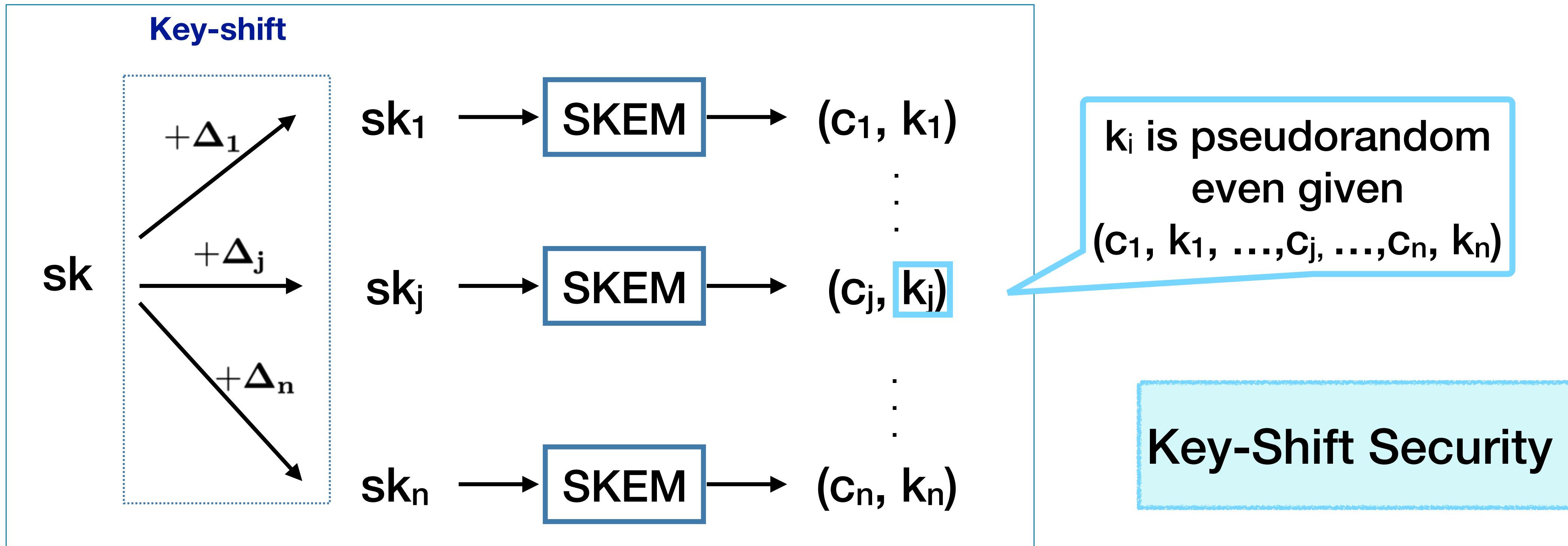
Homomorphic extractor:

$$\text{Ext}(w+w', i) = \text{Ext}(w; i) + \text{Ext}(w'; i).$$

Building Block – SKEM

Symmetric Key Encapsulation Mechanism is similar to traditional KEM.

- $\text{SKEM}.\text{Enc}(\text{pp}, \text{sk}) \rightarrow (\text{c}, \text{k}).$
- $\text{SKEM}.\text{Dec}(\text{c}, \text{sk}) = \text{k}.$



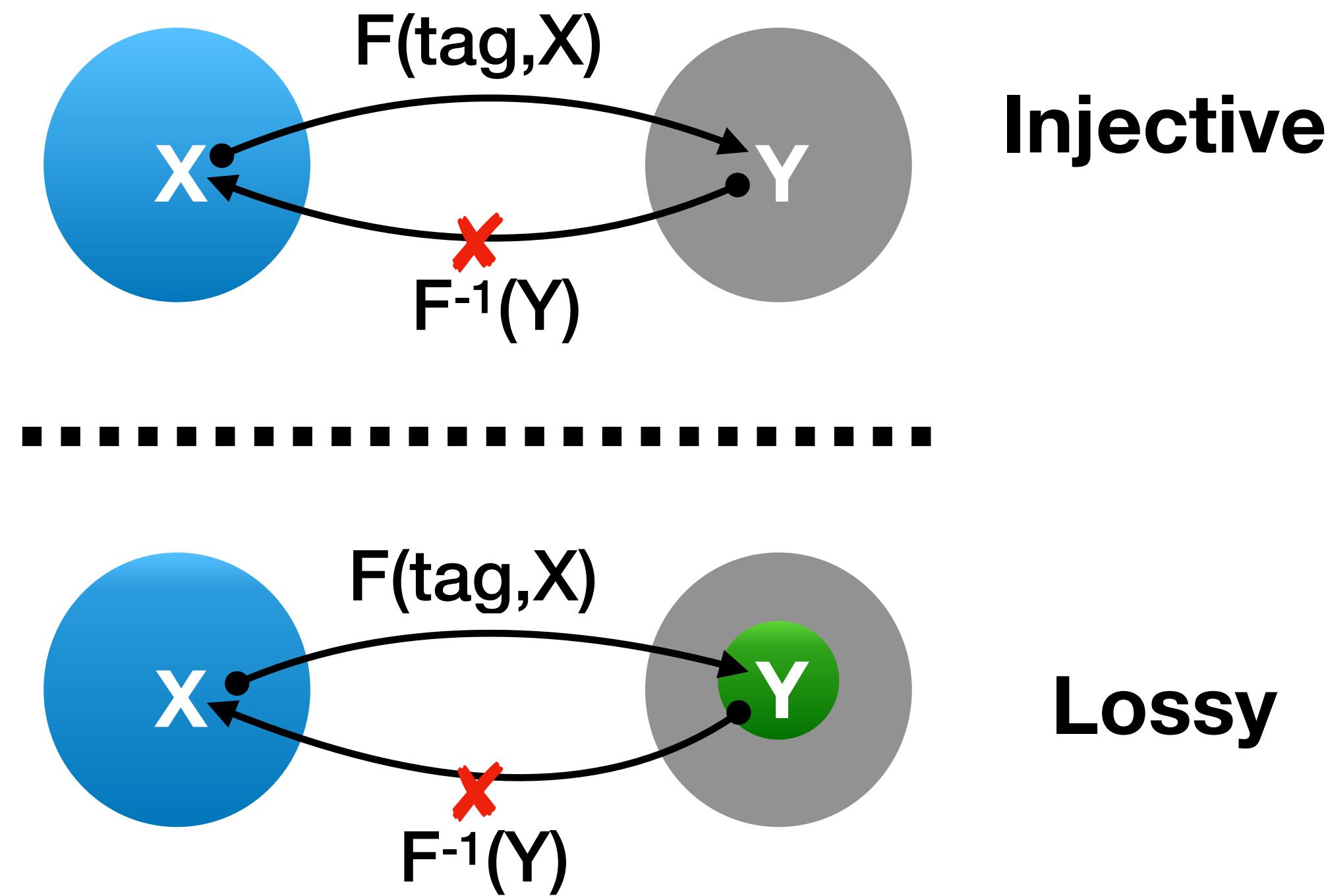
Building Block – LAF

Lossy Algebraic Filter (LAF)

$\text{LAF} : \{F : \mathcal{X} \rightarrow \mathcal{Y}\}$

- $F\text{Gen}(1^\lambda) \rightarrow (F_{\text{pk}}, F_{\text{td}})$.
- $F\text{Eval}(F_{\text{pk}}, \text{tag}, X) \rightarrow Y$.
- $F\text{Tag}(F_{\text{td}}, t) \rightarrow t'$.

$\text{tag} = (t, t')$



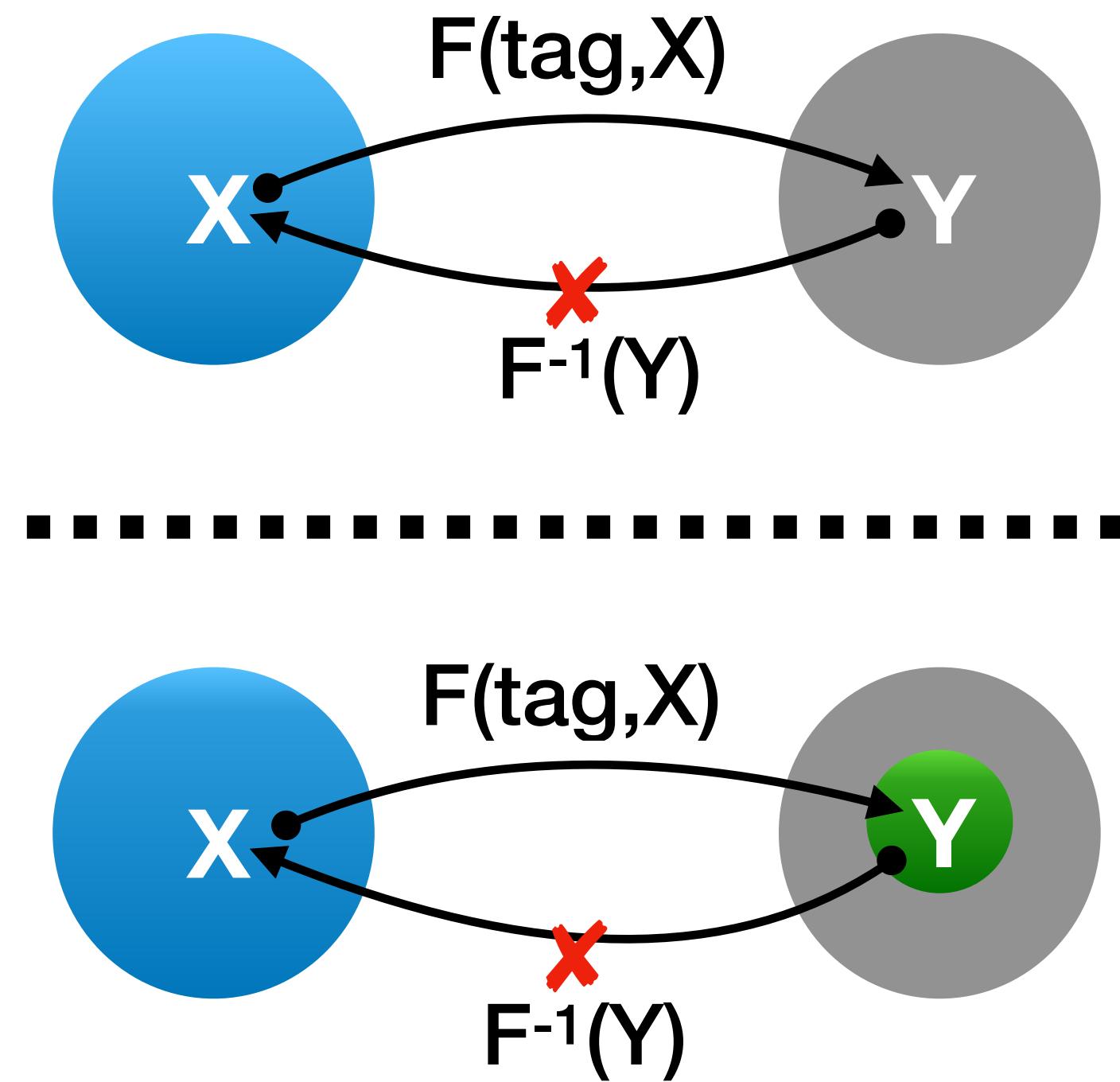
Building Block – LAF

Lossy Algebraic Filter (LAF)

LAF : $\{F : \mathcal{X} \rightarrow \mathcal{Y}\}$

- FGen(1^λ) $\rightarrow (F_{pk}, F_{td})$.
- FEval(F_{pk} , tag, X) $\rightarrow Y$.
- FTag(F_{td} , t) $\rightarrow t'$.

tag=(t, t')



Injective

↔

Lossy

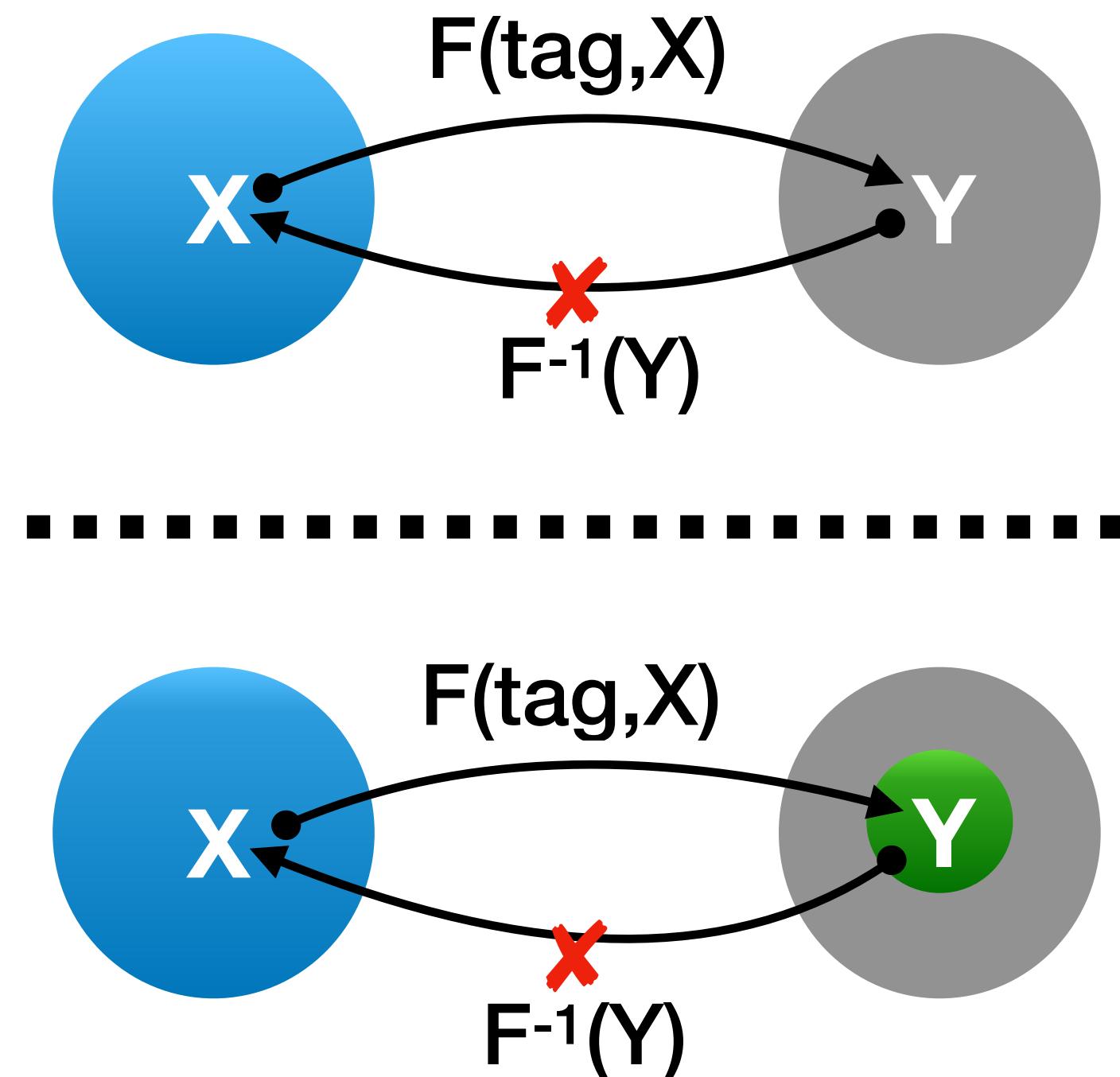
Building Block – LAF

Lossy Algebraic Filter (LAF)

LAF : $\{F : \mathcal{X} \rightarrow \mathcal{Y}\}$

- FGen(1^λ) $\rightarrow (F_{pk}, F_{td})$.
- FEval(F_{pk} , tag, X) $\rightarrow Y$.
- FTag(F_{td} , t) $\rightarrow t'$.

tag=(t, t')



Injective



Lossy

- **Evasiveness:** It is hard to find a non-injective tag without F_{td} .
- **Lossiness:** If the tag is lossy, the function value is only depend on $\sum_{i=1}^n u_i X_i$.

Building Block—LAF

Lossy Algebraic Filter (LAF)

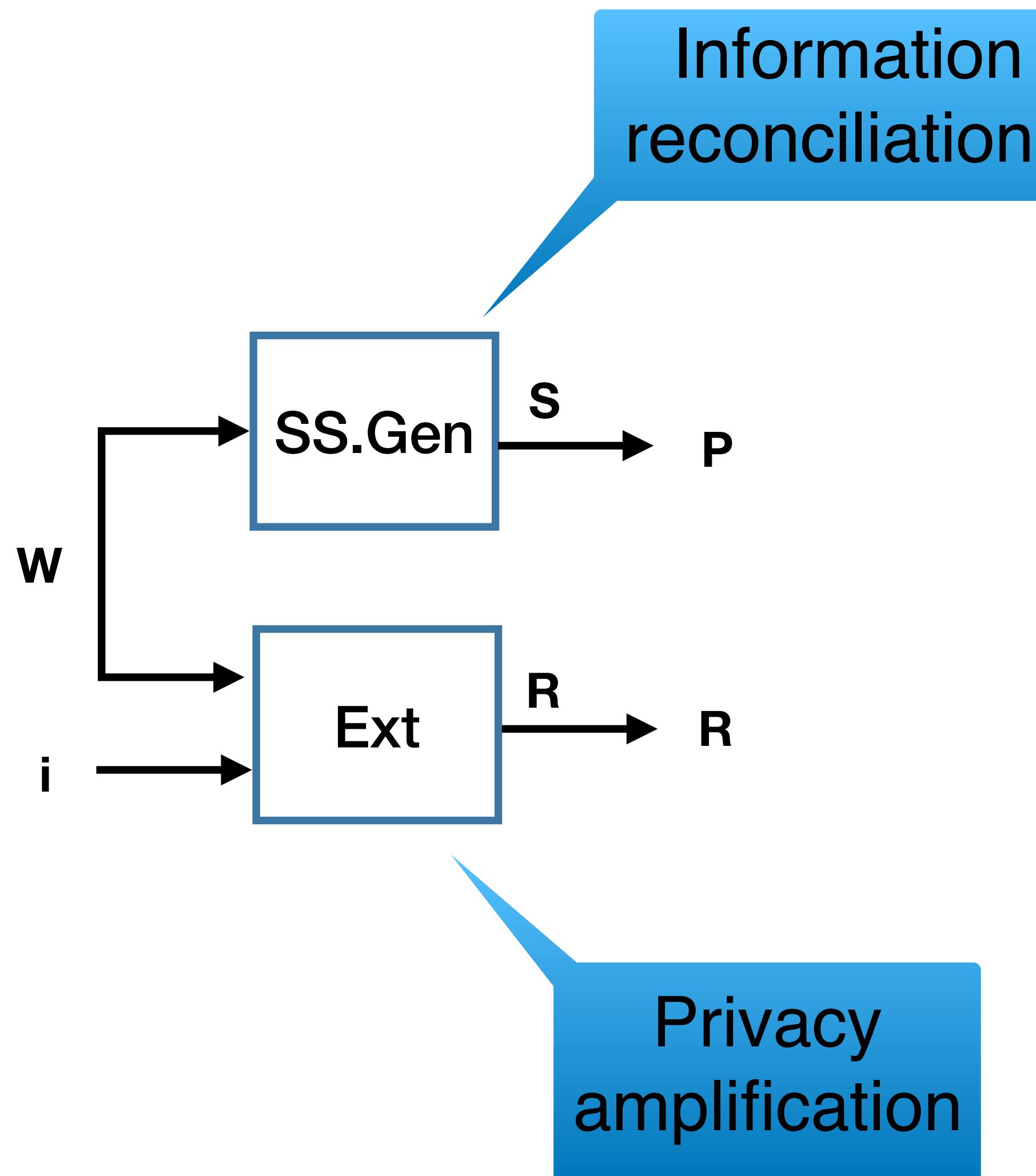
$\text{LAF} : \{F : \mathcal{X} \rightarrow \mathcal{Y}\}$

- $F\text{Gen}(1^\lambda) \rightarrow (F_{\text{pk}}, F_{\text{td}}).$
- $F\text{Eval}(F_{\text{pk}}, \text{tag}, X) \rightarrow Y.$
- $F\text{Tag}(F_{\text{td}}, t) \rightarrow t'.$

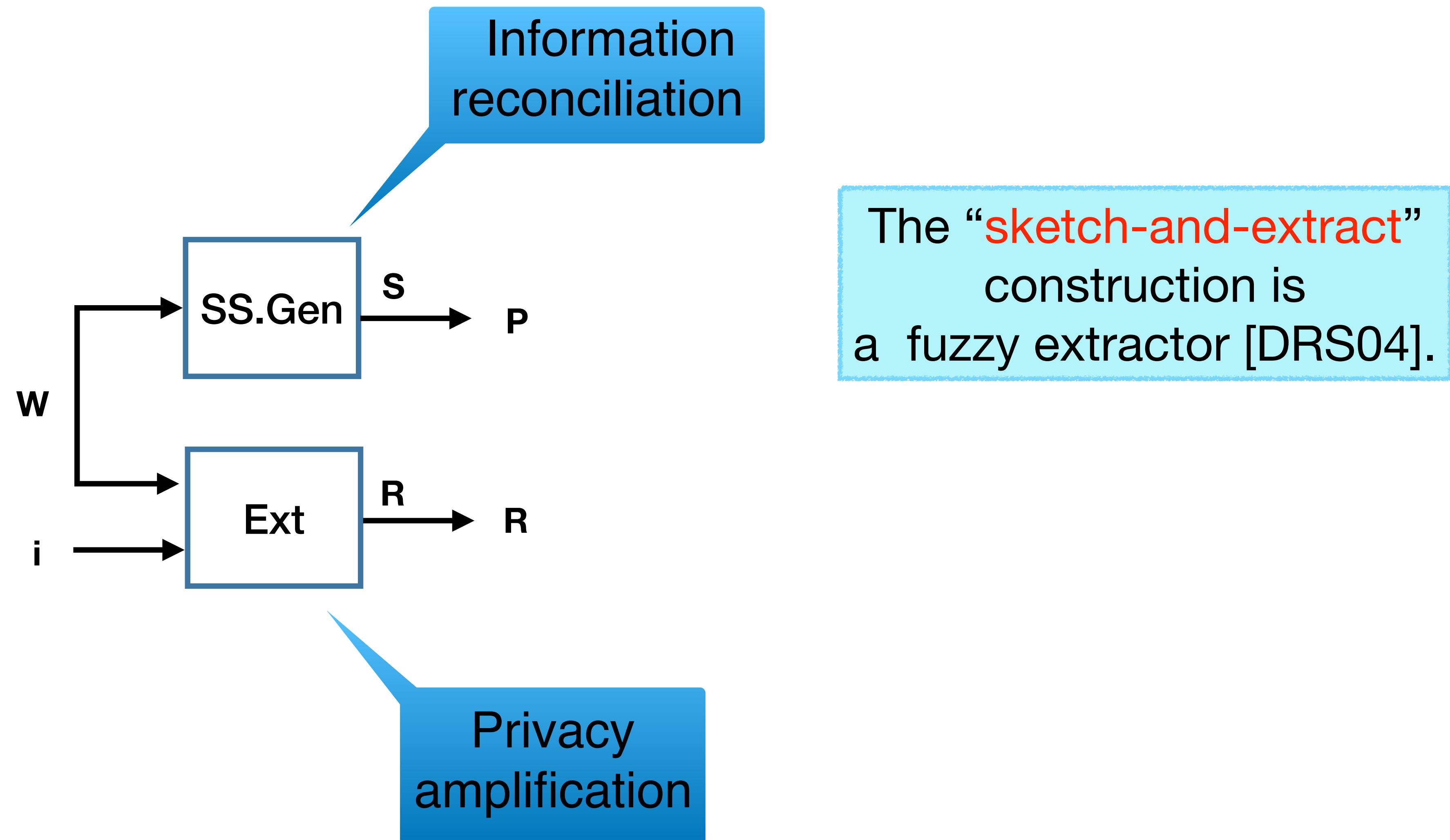
Homomorphic LAF:

$$F\text{Eval}(F_{\text{pk}}, \text{tag}, X_1 + X_2) = F\text{Eval}(F_{\text{pk}}, \text{tag}, X_1) + F\text{Eval}(F_{\text{pk}}, \text{tag}, X_2)$$

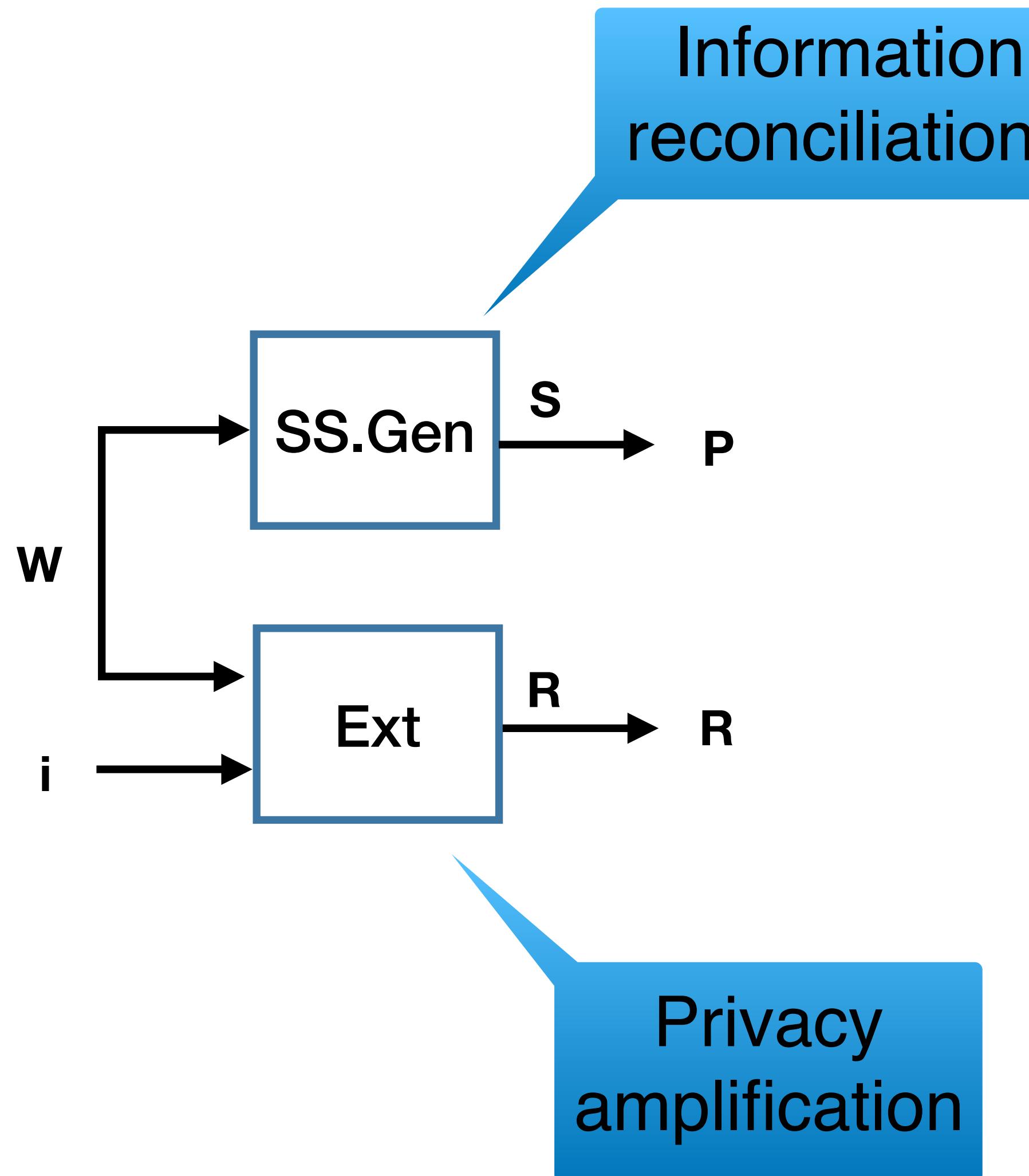
Sketch-and-Extract Paradigm



Sketch-and-Extract Paradigm



Sketch-and-Extract Paradigm

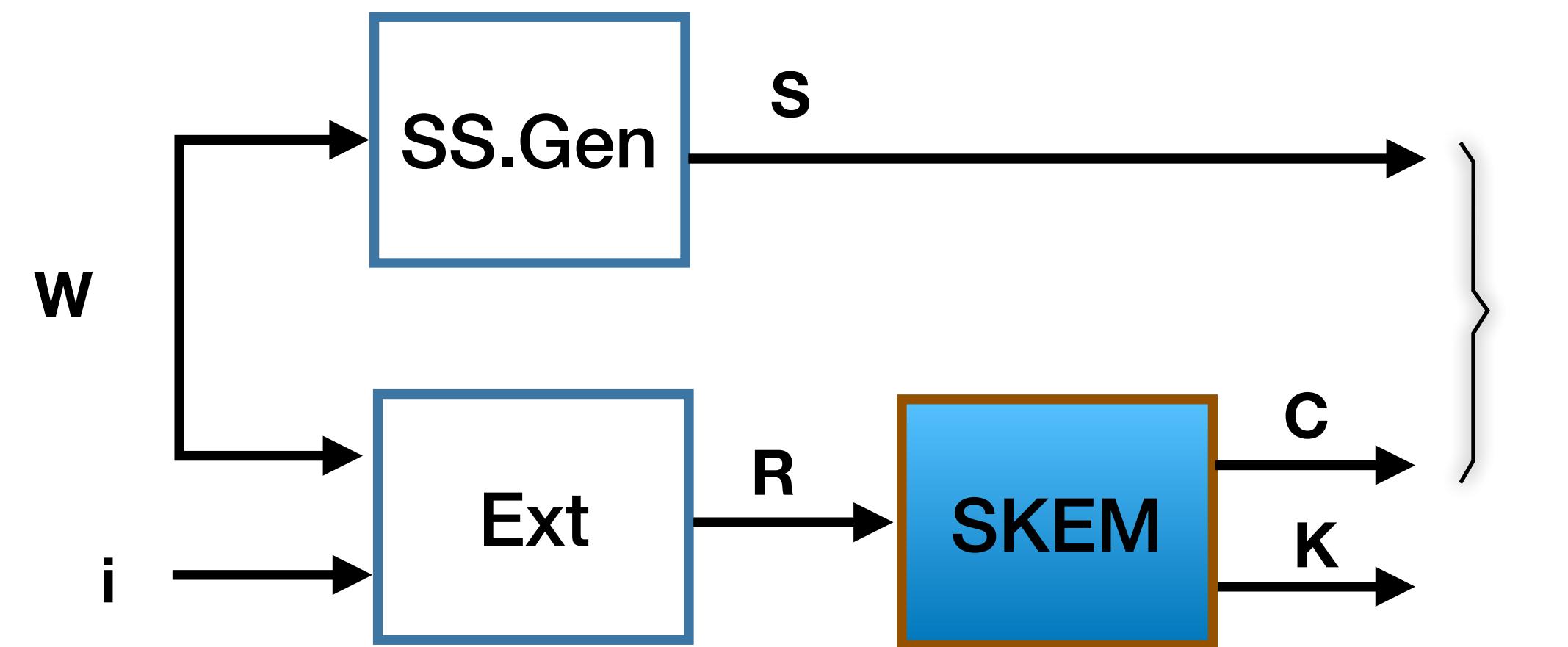


The “**sketch-and-extract**” construction is a fuzzy extractor [DRS04].

Not reusable:
Same w , same R .

Not robust:
No authentication

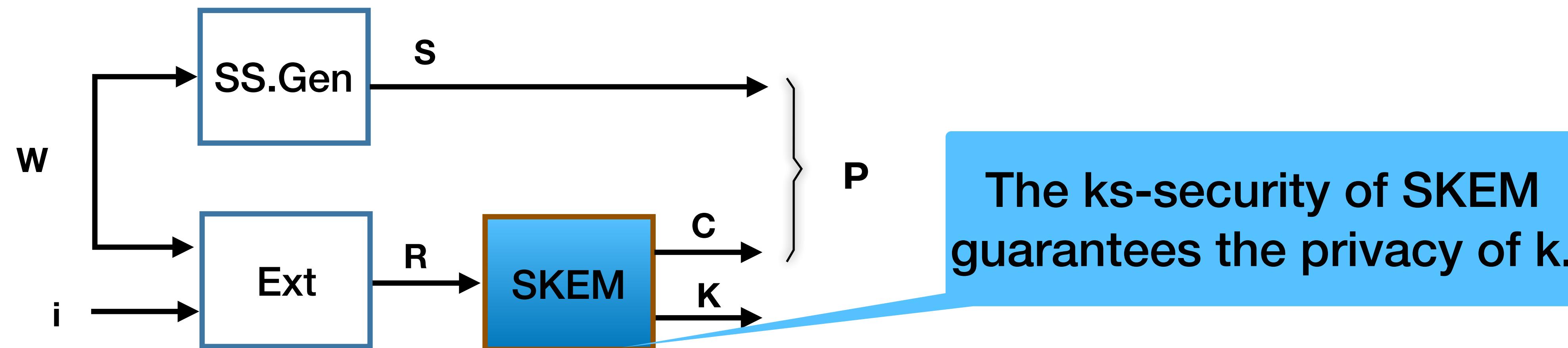
How to Achieve Reusability



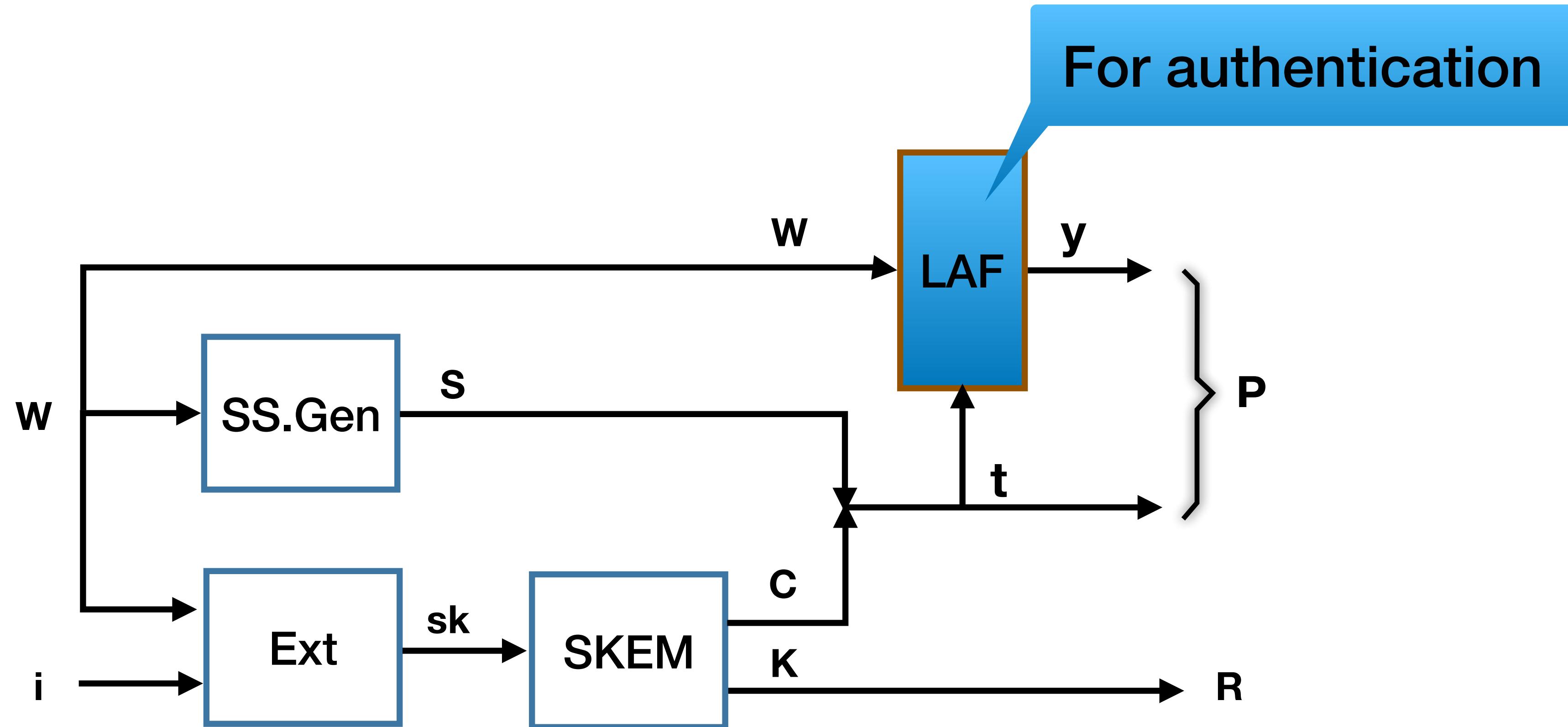
How to Achieve Reusability

Homomorphic properties:

- $s_j = \text{SS.Gen}(w + \delta_j) = \text{SS.Gen}(w) + \text{SS.Gen}(\delta_j)$.
- $R_j = \text{Ext}(w + \delta_j, i) = \text{Ext}(w, i) + \text{Ext}(\delta_j, i)$.



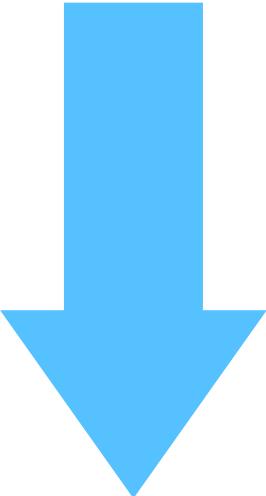
How to Achieve Robustness



Seed i of Ext , Fpk of LAF and pp of SKEM are common reference string, which can be stored publicly, but can not be modified.

Reusability & Robustness

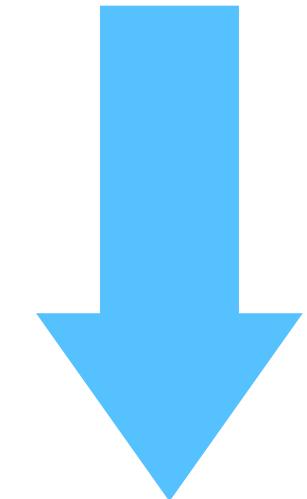
- $s_j = \text{SS.Gen}(w + \delta_j) = \text{SS.Gen}(w) + \text{SS.Gen}(\delta_j).$
- $sk_j = \text{Ext}(w + \delta_j, i) = \text{Ext}(w, i) + \text{Ext}(\delta_j, i).$
- $y_j = \text{FEval}(F_{pk}, \text{tag}, w + \delta_j) = \text{FEval}(F_{pk}, \text{tag}, w) + \text{FEval}(F_{pk}, \text{tag}, \delta_j).$



All tags are changed into lossy tags.

Reusability & Robustness

- $s_j = \text{SS.Gen}(w + \delta_j) = \text{SS.Gen}(w) + \text{SS.Gen}(\delta_j).$
- $sk_j = \text{Ext}(w + \delta_j, i) = \text{Ext}(w, i) + \text{Ext}(\delta_j, i).$
- $y_j = \text{FEval}(F_{pk}, \text{tag}, w + \delta_j) = \text{FEval}(F_{pk}, \text{tag}, w) + \text{FEval}(F_{pk}, \text{tag}, \delta_j).$



All tags are changed into lossy tags.

Enough entropy is left for Ext to extract a key and for LAF to authenticate.

Instantiation

- Homomorphic Ext and SS have information theoretical instantiations.
- Homomorphic LAF can be constructed by the DLIN assumption.
- Key-shift secure SKEM can be constructed by the DDH assumption.

Our rrFE is based on standard assumptions.

Summary

- Our contribution
 - We constructed the first robustly reusable fuzzy extractor from standard assumption.
- Open problem
 - Robustly reusable FE for arbitrary correlated inputs.

Thank you !