

Programming the Demirci-Selçuk Meet-in-the-Middle Attack with Constraints

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Outlines

- 1 Introduction
- 2 Modelling the MITM attack
- 3 Applications in Design
- 4 Conclusion

Outline

- 1 Introduction
 - Description of Demirci-Selçuk MITM
- 2 Modelling the MITM attack
- 3 Applications in Design
- 4 Conclusion

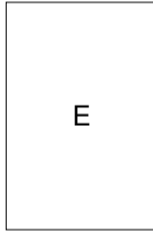
Demirci-Selçuk MITM Attack

- Demirci-Selçuk MITM, FSE 2008 [DS08].
- Various Creative Techniques:
Differential Enumeration, Key Bridging, Key Dependent Sieve, . . . ,
[DKS10, DFJ13, DF13, DF16, LJ16]
- General Model, Dedicated Search Algorithm [LWWZ13, DF13, DF16]

Automatic Searching methods

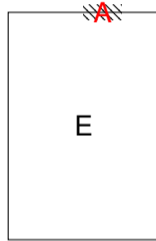
- MILP,CP,SAT,SMT
- Differential, Linear, Integral, 3-subset MITM ...
[KLT15, SHW⁺14, ST17, CJF⁺16, XZBL16, GMS16, Sas18]

MITM Distinguisher



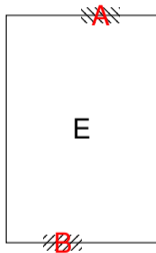
MITM Distinguisher

$\delta(A)$ -set: $\{P^0, P^1, \dots, P^{N-1}\}$



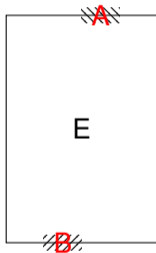
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MITM Distinguisher

$$\delta(A)\text{-set: } \{P^0, P^1, \dots, P^{N-1}\}$$


$$\{C^0, C^1, \dots, C^{N-1}\}$$

$$\Delta_E(A, B): \{C^0[B] \oplus C^1[B], C^0[B] \oplus C^2[B], \dots, C^0[B] \oplus C^{N-1}[B]\}$$

Distinguisher of MITM



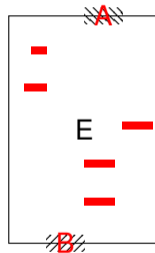
Distinguisher of MITM

- Random: \mathcal{N}_R



Distinguisher of MITM

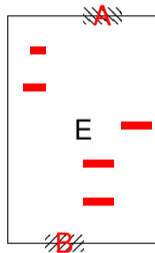
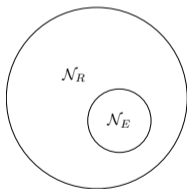
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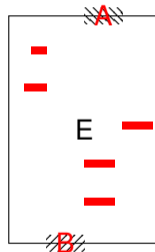
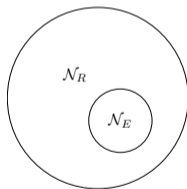
Condition $\mathcal{N}_E < \mathcal{N}_R$



Distinguisher of MITM

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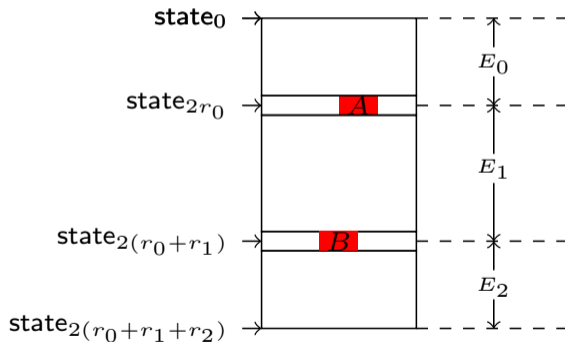
Condition $\mathcal{N}_E < \mathcal{N}_R$



Distinguisher: (A, B, \mathcal{N}_E)

Key Recovery Attack of MITM

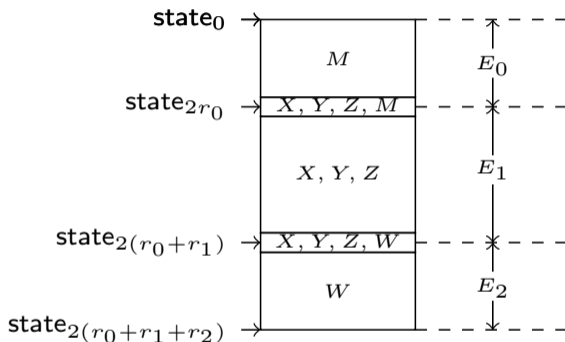
- A cipher is divided in three keyed permutations: E_0, E_1, E_2
- Construct distinguisher (A, B, \mathcal{N}_E) at E_1



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 - Modelling the Distinguisher
 - Modelling the Key-Recovery Process
- 3 Applications in Design
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Variables



- $\text{Var}(X)$ describe the forward differential
- $\text{Var}(Y)$ describe the backward determination
- $\text{Var}(Z)$ model the relation between $\text{Var}(X)$ and $\text{Var}(Y)$

Forward differential

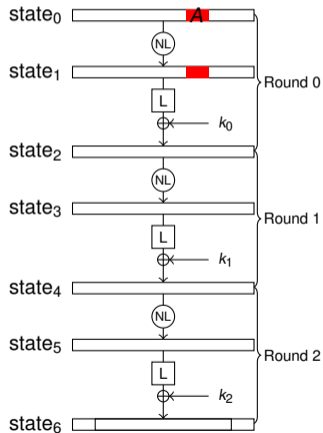
Variables $\text{Var}(X)$

$$X_r[j] = 0 \text{ iff } P_r^0[j] \oplus P_r^i[j] = 0, \forall i \in 1, \dots, N-1.$$

- $\text{Var}(X)$ for state₀

$$X_0[j] = 1 \text{ iff } j \text{ in } A.$$

- X_r propagate to X_{r+1} with probability 1



Forward differential

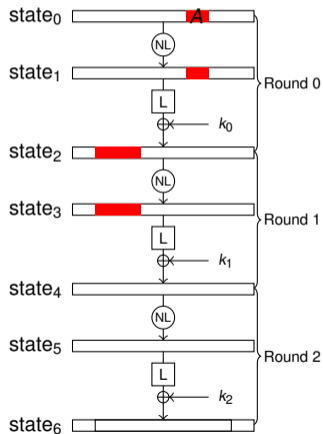
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Forward differential

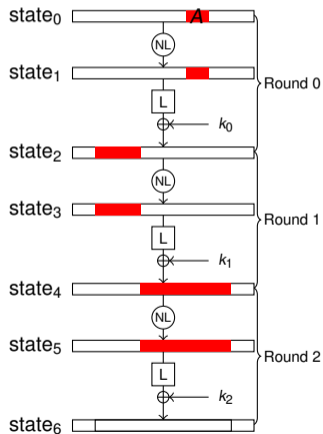
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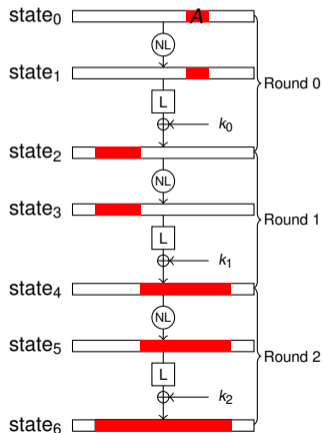
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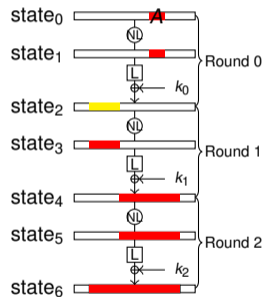
Property of forward differential

values of P^0 at the yellow states

⇓

$$P_6^0 \oplus P_6^i, \forall i \in \{1, 2, \dots, N-1\}$$

$\{P^0[A_{2r}]\}_{r \in 1,2}$ determine $P_6^0 \oplus P_6^i$



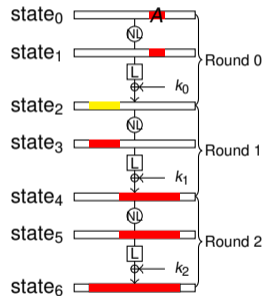
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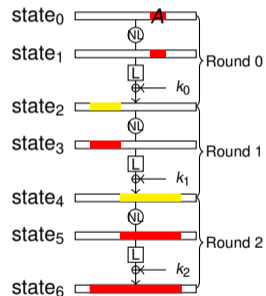
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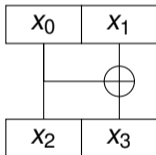
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Forward Differential

Examples for $\text{Var}(X)$

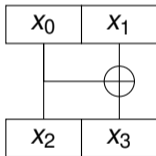
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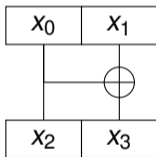
$$2x_3 \geq x_0 + x_1$$

$$x_3 \leq x_0 + x_1$$

Forward Differential

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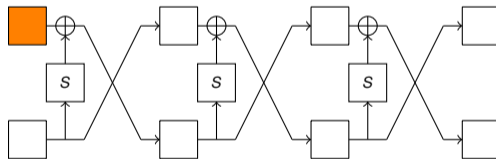
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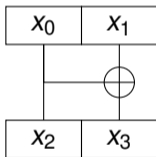
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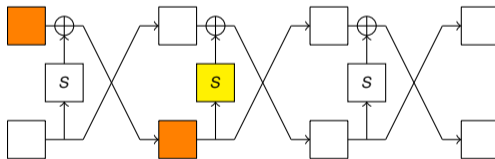
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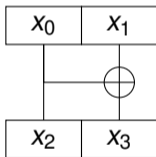
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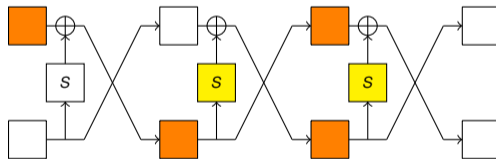
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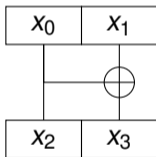
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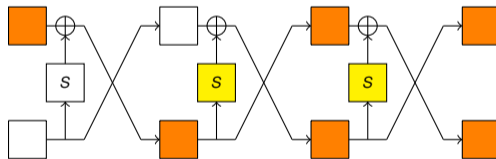
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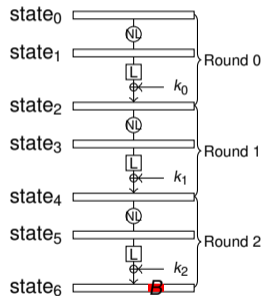
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Backward determination

Variables $\text{Var}(Y)$

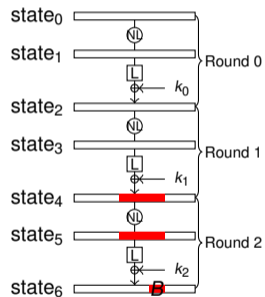
- $\text{Var}(Y)$ for state_6
 $Y[j] = 1$ iff $j \in B$.
- Difference at B are determined by these colored positions



Backward determination

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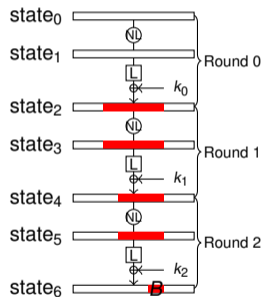
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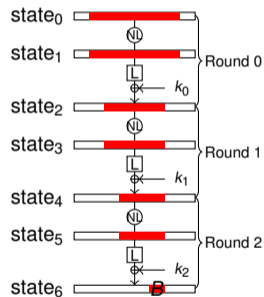
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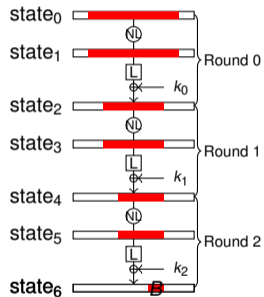
Property of backward determination

Values of P^0 at the yellow nibbles

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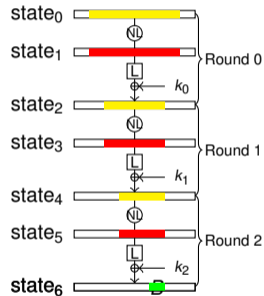
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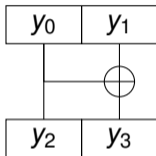
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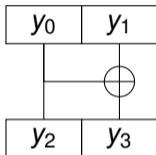
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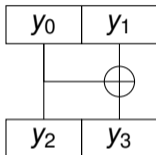
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$$y_2 + y_3 \geq y_0$$

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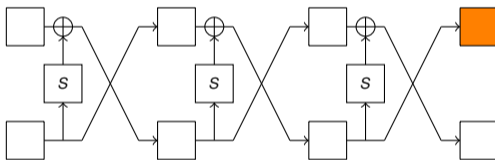
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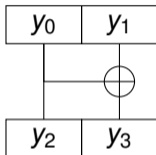
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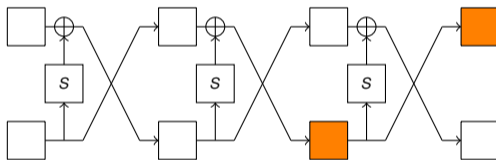
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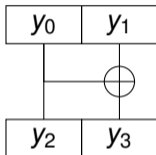
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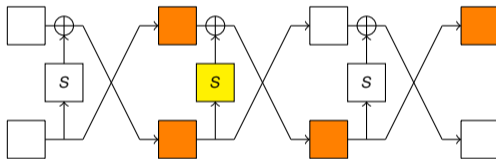
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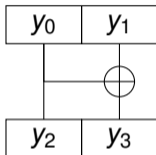
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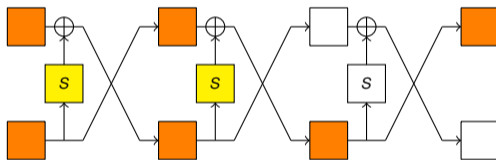
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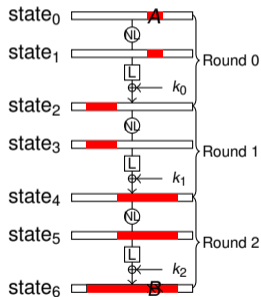
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Constraints for $\text{Var}(Z)$

Variables $\text{Var}(Z)$ describe the relations between $\text{Var}(X)$ and $\text{Var}(Y)$:

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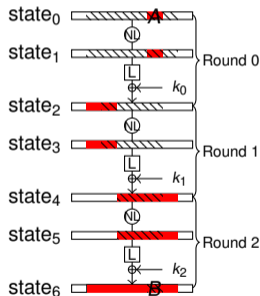


objective function: Minimize $\sum_{r=r_0+1}^{r_0+r_1-1} Z_{2r}$

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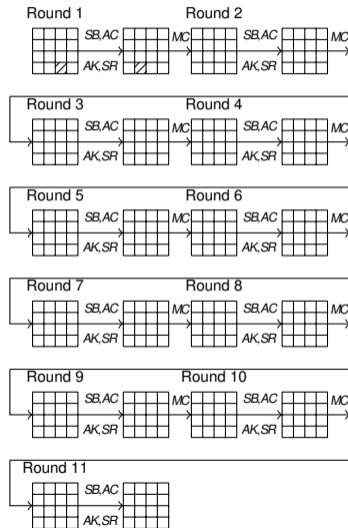
Application to SKINNY

- Proposed at CRYPTO 2016

- $$MC = \begin{pmatrix} 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

$$MC_{i,0}X_r[0] + MC_{i,1}X_r[1] + MC_{i,2}X_r[2] + MC_{i,3}X_r[3] - X_{r+1}[i] \geq 0$$

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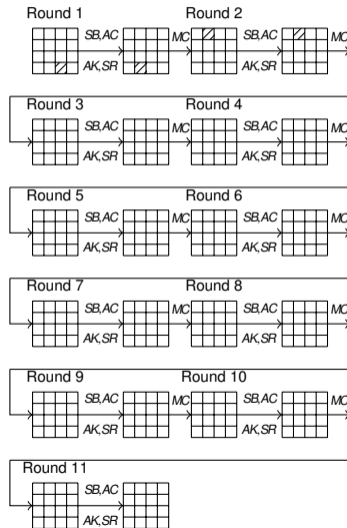
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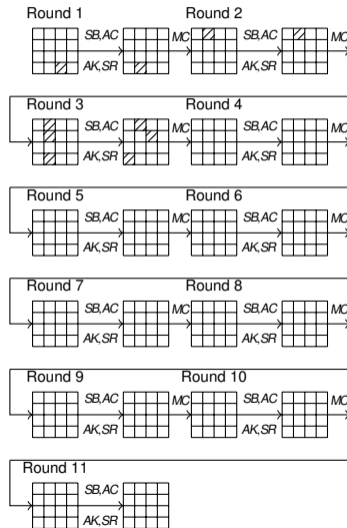
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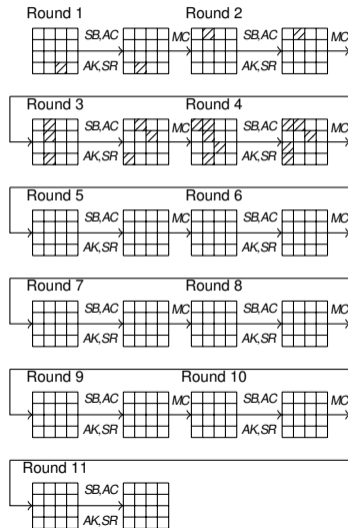
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$$MC_{i,0}X_r[0] + MC_{i,1}X_r[1] + MC_{i,2}X_r[2] + MC_{i,3}X_r[3] - X_{r+1}[i] \geq 0$$

$$4X_{r+1}[i] - MC_{i,0}X_r[0] - MC_{i,1}X_r[1] - MC_{i,2}X_r[2] - MC_{i,3}X_r[3] \geq 0$$



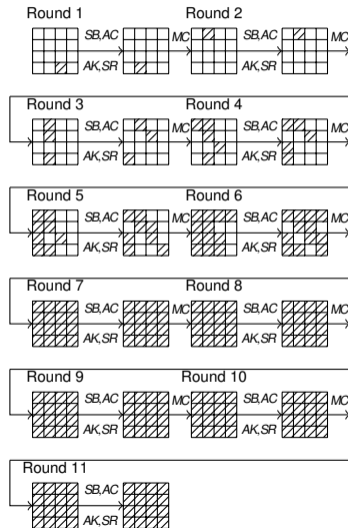
Application to SKINNY

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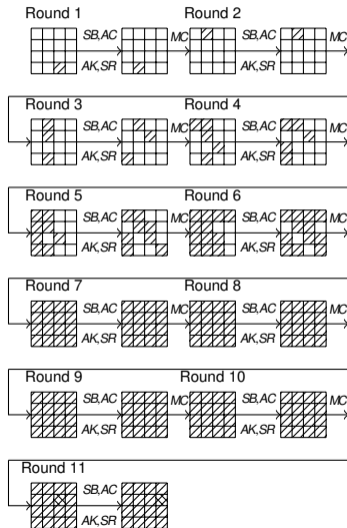
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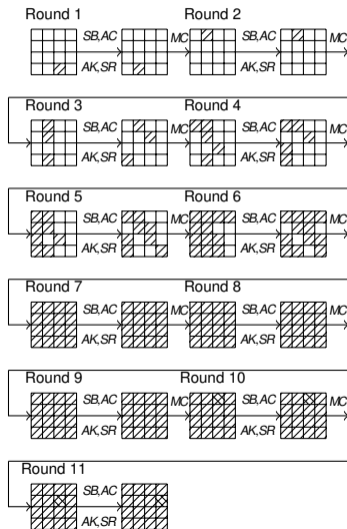
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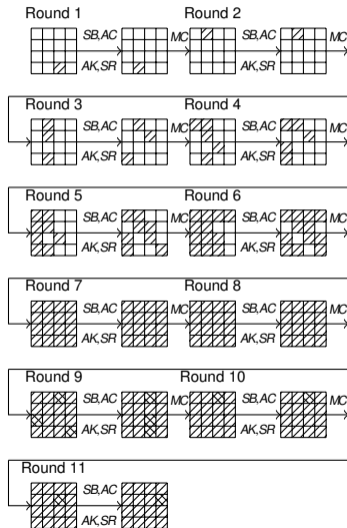
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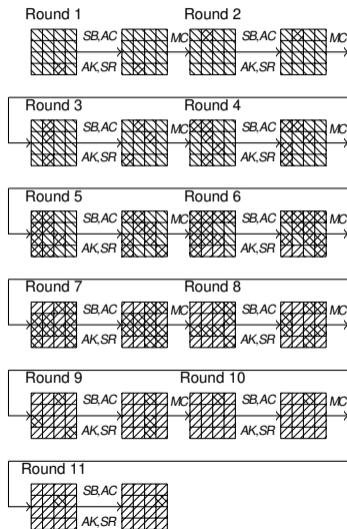
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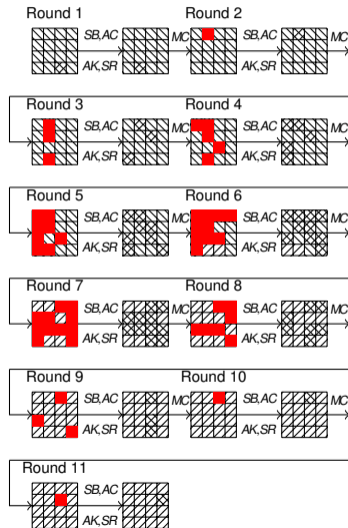
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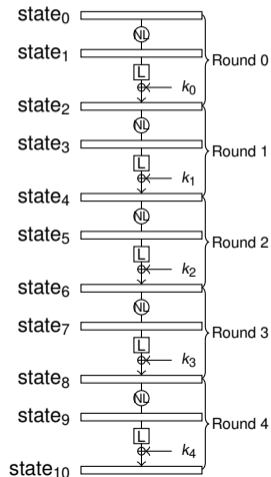
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New 0-1 variables $\text{Var}(M)$ for E_0

$\text{Var}(M)$: Determine the plaintext structure

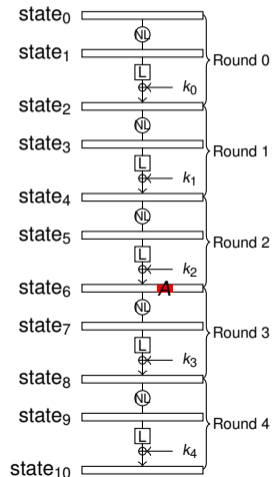
- $M_{2r_0} = X_{2r_0}$
- backward differential propagation with probability 1.
- forward differential through E_0^{-1} :



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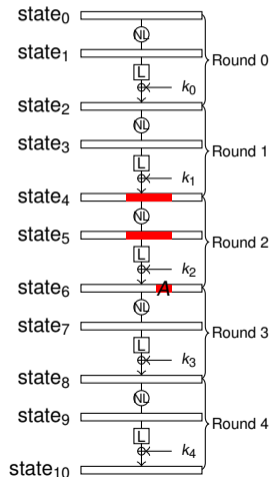
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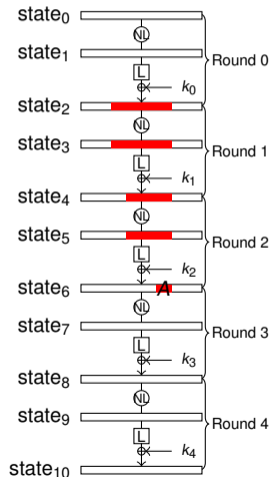
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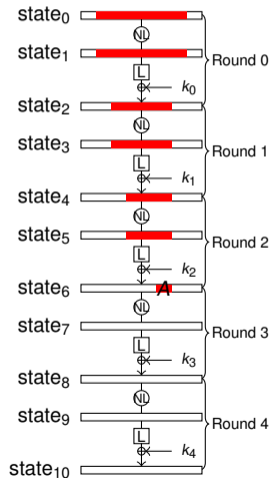
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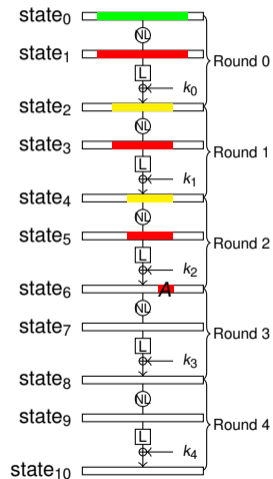
$\text{Var}(M)$: Determine the plaintext structure

- $M_{2r_0} = X_{2r_0}$
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- forward differential through E_0^{-1} :



Property of backward differential

- Guess the values of P^0 at the yellow state
- obtain $\{P^0, P^1, \dots, P^{N-1}\}$ s.t. $\{P_6^0, P_6^1, \dots, P_6^{N-1}\}$ is $\delta(A)$ -set

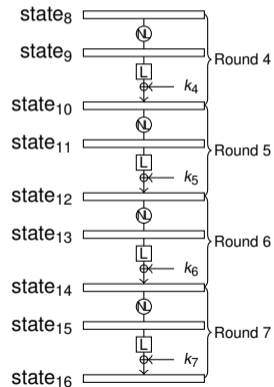


New 0-1 variables $\text{Var}(W)$ for E_2

$\text{Var}(W)$: forward determination process

- $W_{10} = Y_{10}$

- backward determination through E_2^{-1} :

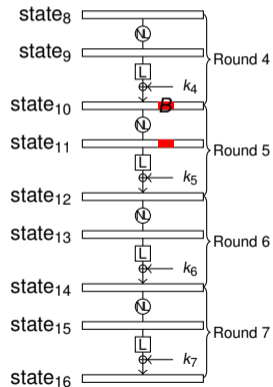


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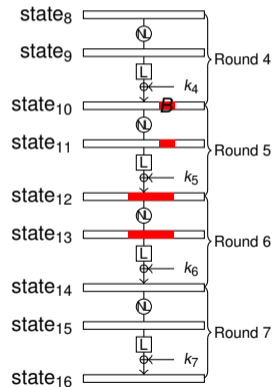


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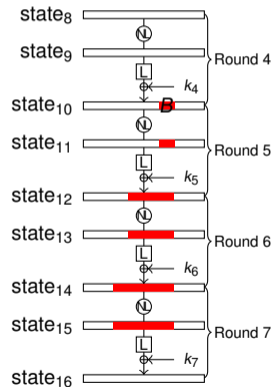


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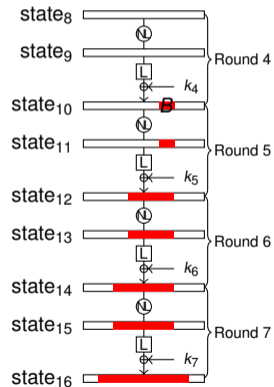


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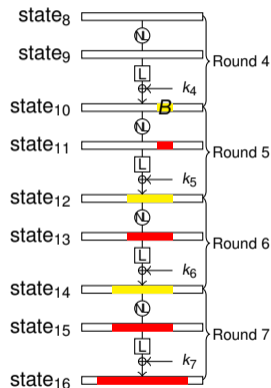


Property of Forward Determination

Guess the values of P^0 at yellow states



Obtain sequence $\Delta_E(A, B)$ at state₁₀

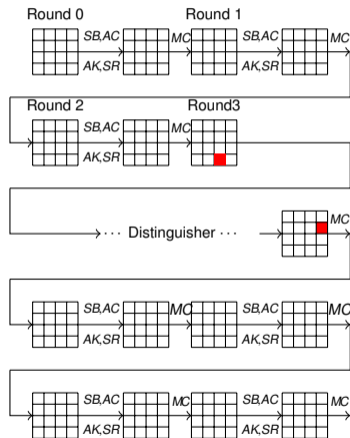


Examples for $\text{Var}(M)$ and $\text{Var}(W)$

$\text{Var}(M)$: Backward differential

$$MC^{-1} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \end{pmatrix}$$

$\text{Var}(W)$: Forward determination

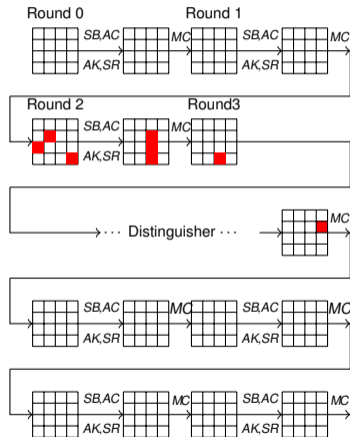


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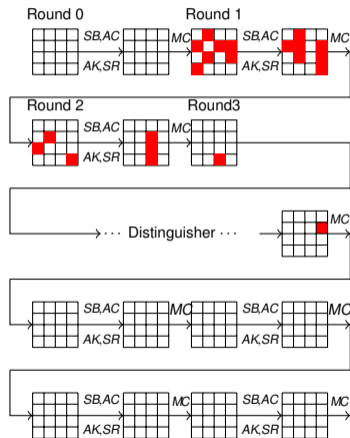


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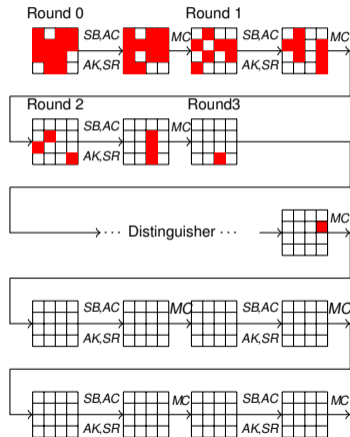


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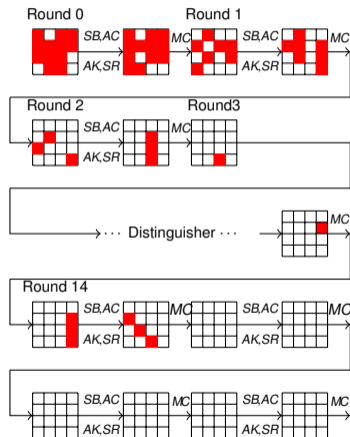


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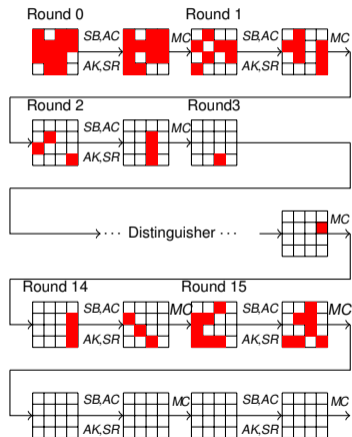


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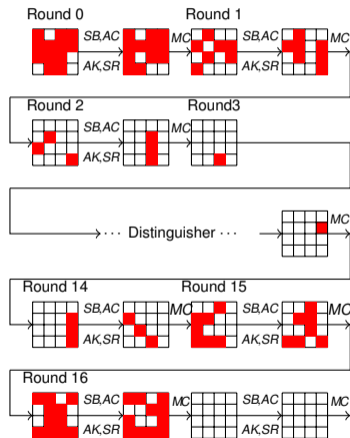


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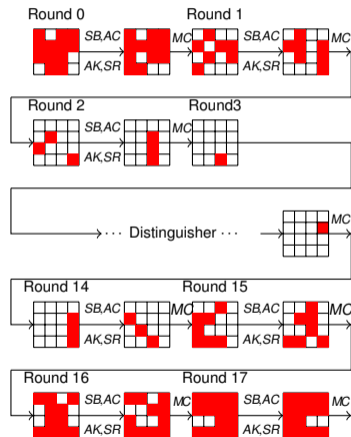


Examples for $\text{Var}(M)$ and $\text{Var}(W)$

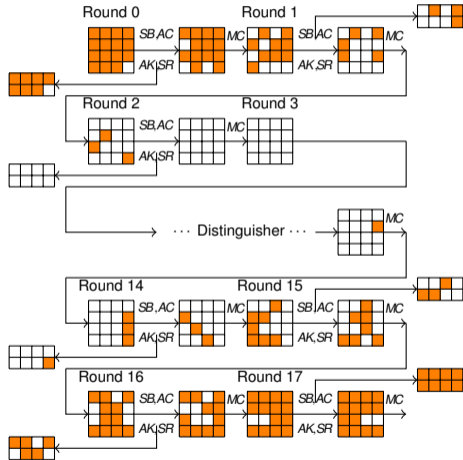
$\text{Var}(M)$: Backward differential

$$MC^{-1} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \end{pmatrix}$$

$\text{Var}(W)$: Forward determination



Guessed Subkeys



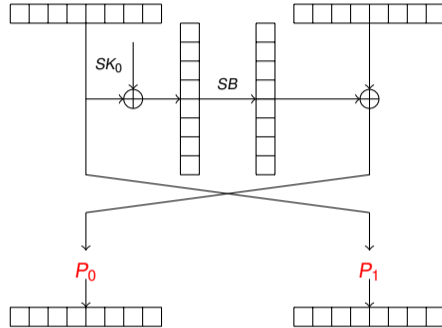
Results

Target	Rounds	Time	Data	Memory	Method	Ref
LBlock	11 + 10	$2^{70.20}$	2^{48} CP	$2^{61.91}$	\mathcal{DS} -MITM	Sect. 7.2
	9 + 0	$2^{74.5}$	—	—	\mathcal{DS} -MITM Dist.	[43]
	23	$2^{74.5}$	$2^{59.5}$ CP	$2^{74.3}$	ID	[47]
	23	$2^{75.36}$	2^{59} CP	2^{74}	ID	[48]
	23	2^{72}	$2^{62.1}$ Kp	2^{60}	MultiD ZC	[47]
	23	2^{76}	$2^{62.1}$ Kp	2^{60}	MultiD ZC	[49]
TWINE80	11 + 9	$2^{77.44}$	2^{32} CP	$2^{82.91}$	\mathcal{DS} -MITM	Sect. 7.3
	23	$2^{79.09}$	$2^{57.85}$ CP	$2^{84.06}$	ID	[50]
	23	2^{73}	$2^{62.1}$ KP	2^{60}	MultiD ZC	[47]
TWINE128	11 + 14	$2^{124.7}$	2^{48} CP	2^{109}	\mathcal{DS} -MITM*	[34]
	25	$2^{124.5}$	$2^{59.1}$ CP	$2^{78.1}$	ID	[34]
	25	2^{119}	$2^{62.1}$ KP	2^{60}	MultiD ZC	[47]
	25	$2^{122.12}$	$2^{62.1}$ KP	2^{60}	MultiD ZC	[49]
SKINNY-128-384	10.5 + 11.5	$2^{382.46}$	2^{96} CP	$2^{330.99}$	\mathcal{DS} -MITM	Sect. 7.1
	11 + 11	$2^{373.48}$	$2^{92.22}$ CP	$2^{147.22}$	ID	[51]

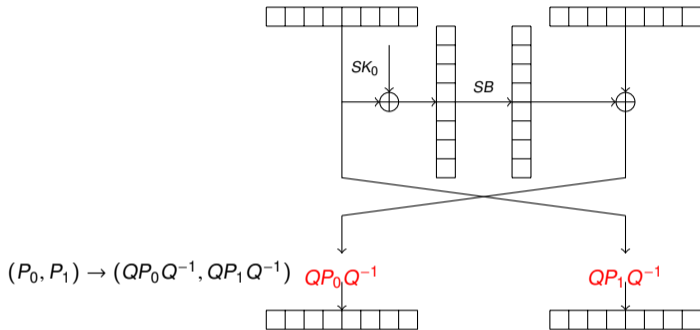
★ We find the attacks with the same complexity.

Outline

- 1 Introduction
- 2 Modelling the MITM attack
- 3 Applications in Design**
 - Results of TWINE and LBlock Structure
- 4 Conclusion



Enumeration: $8! \cdot 8!$



Enumeration: $8! \cdot 8! \rightarrow 22 \cdot 8!$

Results

- 144 permutations: no 15-round IM Distinguisher.
- 12 permutations may be best: no 11-round MITM distinguisher

Outline

- 1 Introduction
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Conclusion

Conclusion

- modelling the MITM attack
- IM and MITM for variants cipher of LBlock and TWINE

Future Work

- Differential Enumeration
- Key Bridging
- ...

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Thanks for your attention.